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ON-LINE PROGRAM ANALYSIS
AND DIAGNOSIS

JOHN L. LIGHTSTONE

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ON-LINE PROGRAM ANALYSIS
AND DIAGNOSIS

by

John L. Lightstone

Lieutenant, United States Navy

Submitted in partial fulfillment of
the requirements for the degree of

MASTER OF SCIENCE
IN
ENGINEERING ELECTRONICS

United States Naval Postgraduate School
Monterey, California

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ABSTRACT

For many program errors and program checkout problems, on-line techniques provide a promising method of attack. An approach developed in connection with time-sharing computer systems is examined. Then Program Trace, an on-line diagnostic program developed by the author for the CDC 1604 computer, and a Data Display Model DD 65 display and control console is presented and examined in detail.

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ABBREVIATIONS

A	Arithmetic Register
ADD	Execution Address
E^i	Index Register Number i
EBN	Bolt, Beranek and Newman, Inc.
BCD	Binary Coded Decimal
CDC	Control Data Corporation
DD	Data Display, Inc.
FC	Function Code
P	Program Address Register
Q	Auxiliary Arithmetic Register
SDC	Systems Development Corporation

1. Introduction.

Much has been done to automate the error detection processes for program checkout. On the other hand, it is apparent that a serious problem still exists. A significant portion of the time spent in the development of programs, both in terms of man hours and in terms of computer time, is in the checkout phase. Analysis of certain types of checkout problems is not easily automated. New approaches to the diagnostic problem are needed.

One promising approach is the use of the computer to provide on-line assistance to the programmer in analyzing his programs. In order to investigate the potentials and problems of this approach, an on-line diagnostic program was developed for use with the CDC 1604 computer and the DD 65 display console at the Naval Postgraduate School. Because of the technique used in the central part of the program, it is called Program Trace.

Before looking at Program Trace, it will be useful to catalog the types of checkout problems and program errors that occur and to list the aids and techniques available for attacking these problems, in order to determine the weak areas. Then a look will be taken at some on-line diagnostic programs, including a detailed examination of Program Trace.

2. Checkout Problems and Program Errors.

The checkout problems and program errors will be considered in two phases, those that occur or appear prior to or during compilation, and those that appear during the running of the program. Those in the first phase will be broken down into

clerical and syntactic errors, programming errors, and machine failures. Those in the running phase will be broken down into programming errors, conditional errors, and machine failures.

2.1 Pre-compilation and Compilation Phase.

The clerical and syntactic errors are the typographical mistakes, mispunched cards, and format errors which result in meaningless symbols or instructions.

Programming errors which appear during compilation are mistakes such as exceeding the storage capacity of the machine by reserving too much space for arrays or other data storage; not defining symbolic addresses; not identifying a reserved array; and calling for sub-programs not available on the system library.

Machine failures include anything from the dropping of a bit during transfer to or from magnetic tape to the failure of a logic component in the machine to the failure of a power supply. Also, included are failures of external equipment associated with the computer.

2.2 Running Phase.

Machine failures have no reference to the phase of the program. They are the same problems as already mentioned.

Conditional errors are those errors which depend on the status of certain parameters in the program or on the intermediate results of the computation, and are not always foreseeable. Examples are calculations resulting in a negative

argument for a square root function and the overflow of an arithmetic register during calculations.

Programming errors are flaws of a fundamental nature in the logical structure of the program and include a variety of problems. Three of these are the occurrence of unexpected jumps, deadends in loops, and programs that run to conclusion but with improper results.

The unexpected jumps are of two types - those that jump to an unused memory cell and stop the computer, and those that jump to another part of the object program or into the system programs and do not stop the computer immediately. In either case, it is difficult to determine where the jump originated.

Unending loops can occur from unindexed jumps or from conditional jumps for which the condition is never satisfied. In a dual instruction machine, such as the CDC 1604, a dead-end can also occur with a skip instruction in the lower half of a memory cell making continuous half exits on itself.

Programs which run to conclusion but with improper results are listed here not because they are a type of error of themselves, but because, quite often, this is the only external indication of internal programming errors. The cause can be an unexpected jump, as already mentioned, or a flaw in the logic used by the programmer in the construction of the program.

Other examples of programming errors are not supplying the required arguments for a sub-program, having input data for the program in improper format, and indexing a variable

array out of the storage area reserved for it in memory.

3. Diagnostic Aids and Methods.

The aids available during the pre-compilation and compilation phase include the error print-outs of the compiler and of the system executive program. The former giving indications of typographical and format errors and the latter giving indications of illegal procedures and some machine failures. These print-outs are well developed and will detect and pinpoint most of the clerical and syntactic errors.

Another valuable aid is a machine language or symbolic machine language listing of the object program which can be supplied by the compiler during compilation.

The diagnostic aids available during the running phase are generally limited to error print-outs by system library sub-programs when supplied with illegal arguments, and to console indications of the contents of the operational registers when the computer stops, console indications of registers overflowing, and console indications of hang-ups during input or output operations.

The diagnostic methods used during and after the running phase are generally variations of four techniques: The use of extra Print statements, Core Dumps, breakpoints, and stepping through the program.

By using extra Print statements at selected points in the object program, the value of parameters in the program can be determined at these points and an indication can be obtained of where trouble occurs.

The Core Dump is used to provide a printed record of the contents of the computer memory or selected portions of it for off-line analysis. This may be done at the completion of the run or at some intermediate point and in either a numeric or a mnemonic format.

The use of breakpoints to stop the computer at selected locations in the object program is valuable for examining and changing conditions at these points. However, it is not always known where it is desirable to stop.

In some cases where the trouble is localized or the action of a particular portion of the program is not understood, stepping through the program one instruction at a time is useful in order to follow the program execution exactly.

4. Weak Areas.

As mentioned, the diagnostic aids available during the compilation phase are well developed and are effective in detecting most of the errors which occur during that phase.

On the other hand, the diagnostic methods available in the running phase, while very useful, are often ineffective for determining the causes of some conditional errors or for locating errors in the logical structure of a program. As

a result, it is often necessary to use these methods in a series of repeated runs. The program is run; the programmer studies the results and makes changes and corrections; and another run is made, with the cycle repeated, perhaps many times until the errors are eliminated.

In a situation such as this, full use is not made either of the computer's potential or of the programmer's time, primarily because of the lack of communications between the computer and the programmer while the program is running. The computer does not know the programmer's intentions if they are not explicitly written into the program, and at other times, the programmer is limited in his communications with the computer.

5. On-line Diagnostic Systems.

The major work in developing on-line diagnostic programs has been in conjunction with the development of time-sharing computer systems. This is a natural result, since a major drawback of on-line systems is the inefficient use of computer time. With the ability provided by the time-sharing systems for several people to work on-line simultaneously, this drawback is eliminated.

Two time-sharing systems which have incorporated on-line diagnostic programs into the systems are those of Bolt, Berenac, and Newman [1,2] and the Systems Development Corporation [3]. The BBN system uses a Digital Equipment

Corporation PDP-1 computer and the SDC system uses an FSQ-32 computer with a PDP-1 for an input-output computer.

The operation of the two diagnostic programs is quite similar. Each operator who is running a program has a typewriter keyboard with which he can communicate with either his own program or with the time-sharing executive program and the diagnostic program (through the executive program). A prefix attached to the typed input indicates the program for which it is intended.

The two basic capabilities of the programs are the setting of breakpoints and the ability to examine and change the contents of memory cells, with variations available in how these are used.

Breakpoints (up to three at any one time) may be set by typing a short one or two character command followed by the breakpoint address. The break is accomplished by inserting a jump instruction at the desired break address which will jump into the diagnostic program. When the breakpoint is reached, the diagnostic program types a short message, including any information that may have been previously requested, on the typewriter, and waits for further instructions from the programmer.

Contents of memory cells, within the limits of memory storage used by the program, may be examined by typing another short command followed by the desired address.

Sections of memory of up to twenty cells or the operational registers may also be selected in this manner. Changes may be made in selected cells or registers by the use of another command followed by the new desired contents.

Breakpoints may be set or cells examined before the program is started, after it is finished, while it is stopped at a breakpoint, or while the program is running. In addition, through the executive program, the object program may be started, stopped, restarted, or started over again from the beginning.

The feature that makes these diagnostic programs so useful is their accessibility and ease of use. They are always available to all users of the time-sharing system without the use of any special calling procedures. The diagnostic requests may be typed at any time, and they are sent directly to the diagnostic program by the system executive program.

There are limitations, however. In the case of problems such as unexpected jumps, there is no method for tracing or following the execution of the program as it is running.

A possible means of providing a tracing capability in these systems is by the following method. Provide an additional function in the diagnostic routine which would make a search of the object program before it is run. As a result of the search, the execution address of each jump instruction

in the object program would be changed to that of a position in a table which would be constructed in a reserved area at the end of the program. The table could be designed so that upon each jump into the table, the number of times that jump has been executed and/or the order in which the jumps have been executed could be recorded before a jump is made back to the originally assigned address in the object program.

There would be some problems in implementing this function, and it would make additional time and memory space requirements on the computer system, but it would be very useful for some problems. A special calling procedure would have to be provided for the jump table function, since it could not be as freely accessible as the rest of the diagnostic program.

Program Trace is a program, written by the author, designed to provide on-line diagnostic functions, including a tracing function in a non time-sharing system. However, before discussing Program Trace, a description of the facilities will be provided.

6. Computer Facilities.

The computer facilities at the Naval Postgraduate School include a CDC 1604 computer [6] and a CDC 160 computer in the school's computer center and a Data Display model DD 65 display console [8] and another CDC 160 computer in the Digital Control Laboratory.

The CDC 1604 is a 48 bit, dual instruction computer with a 32K core memory. It has one high speed unbuffered input-output channel and three buffered input and three buffered output channels.

The Fortran 60 System [7] is the Fortran processor in use with the CDC 1604. When the processor is loaded in memory, the resident program occupies the first 4000₈ cells, and the compiler occupies the next 14000₈ cells, so that as each object program is loaded into memory, it begins at address 20000₈. Any library sub-programs which are called are loaded at the end of the object program. The Fortran compiler is designed for either normal Fortran coding or symbolic machine language coding, or for intermixing of the two.

The two CDC 160 computers are 12 bit, single address machines, with 4K core memories, of which only the first 64 cells are directly addressable. They each have one input and one output channel.

The DD 65 is a display and control console, especially configured by Data Display for the Naval Postgraduate School. The console is shown in Figure 1. The associated core memory and logic are in a separate cabinet.

The console has two 12" cathode ray tubes for display of alpha-numeric and vector symbols. For other applications, radar video may also be displayed on the left tube. Display

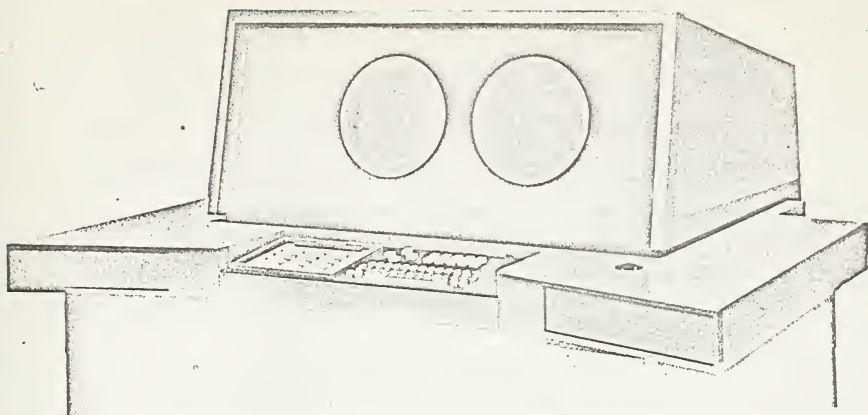


Figure 1. DD 65 Console

commands are stored in the 512 word addressable core memory. The commands are 48 bit words which each correspond to one CDC 1604 word or four CDC 160 words.

The contents of the memory are scanned and displayed approximately 20 times a second to provide a flicker free presentation. The memory can be updated by either the CDC 1604 or the CDC 160 (but not by both at the present time) at any time during the scan. A complete read-write memory cycle is 12.8 micro-seconds.

The position of a character on the screen is designated in X-Y coordinates with the origin at the center of the screen.

As shown in Figure 2, the X (Y) coordinate is given as a nine bit number which increases from 000 at the origin to 377 at the right (top) edge, and from 400 at the left (bottom) edge to 777 at the origin. After an initial position is designated by the use of a designator word, successive characters may be incremented either horizontally or vertically from that position by the use of string words. The formats for designator and string words are shown in Appendix I.

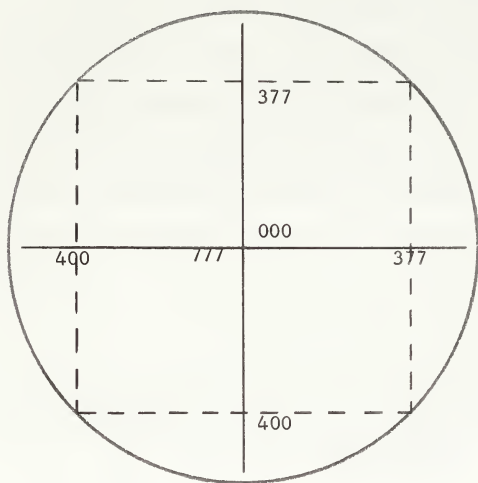


Figure 2. Display Tube Coordinates

The keyboards, Figure 3, are used to provide outputs to the computer. Since there is no direct connection with the display generation circuits or display memory, all control and display functions are performed through the computer. The typewriter keyboard is referred to as keyboard #1, and all other keys and buttons are referred to as keyboard #2.

Cables and a switch on the DD 65 memory cabinet are provided so that the display may be operated with either the CDC 1604 or with the CDC 160 in the Digital Control Laboratory. In the latter case, the CDC 160 may be operated independently or as a satellite of the CDC 1604. For satellite operations, communications between the CDC 160 and the CDC 1604 are through one set of the CDC 1604's buffered input-output channels. When the DD 65 is operated directly with the CDC 1604, communications are via the high speed input-output channel of the CDC 1604.

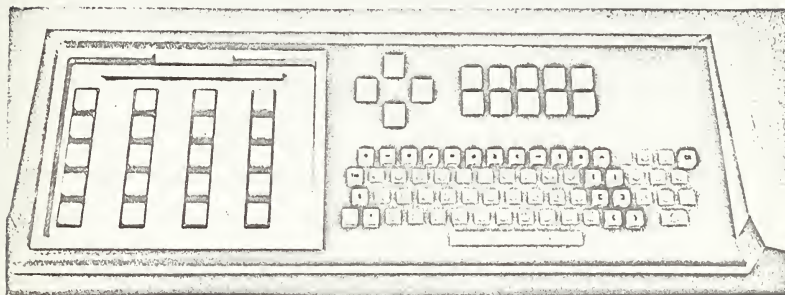


Figure 3. Keyboards #1 and #2.

In considering which configuration to use in implementing Program Trace, both the direct connection of the display to the CDC 1604 and the satellite arrangement have advantages. In the latter case, a load is taken off of the CDC 1604 by having the display generation programs located in the CDC 160. Also, this configuration might more easily be combined with other systems in the future. On the other hand, the direct connection is easier to implement initially. For this reason, the direct connection was chosen.

7. Operation of Program Trace.

Program Trace is written in the form of a sub-program. It is designed so that it can be made a part of the resident library and be callable by the object program (the program to be diagnosed). At present, it is loaded at the end of the object program.

Program Trace is intended for the analysis of machine coded programs. However, it can be useful for diagnosis of Fortran coded programs if a machine code listing of the program is obtained. The Fortran compiler can provide such a listing.

One argument must be provided when Program Trace is called, and that is the address of the instruction immediately following the calling statement. This is done by the use of two symbolic statements prior to the Fortran calling statement.

ENA (L+3)
STA (ADDRESS)
CALL TRACE (ADDRESS)

These statements are normally placed near the beginning of the object program, although they may be placed at other positions if only the latter portion of the program is to be traced.

When the program is run, the Trace program takes control of the execution and first checks to see if the DD 65 is available and ready. If not, the object program is run to completion under the control of Trace, and a table is prepared of each jump instruction executed in the object program including the address of the jump instruction, the address to which it jumped, and the number of times it was executed. The table is printed out in blocks of 50 jump instructions as the table is filled or at the end of the program.

If the DD 65 is available and ready, then its memory will be erased, its keyboard lights set to correspond with initial Trace control flag settings, and an introduction will be displayed on the left tube. The introduction will give the location of the object program in memory, and by giving the location of Trace, will indicate where the object program ends.

PREAMBLE OF YOUR PROGRAM BEGINS AT CELL	20000
INSTRUCTION AFTER 'CALL TRACE' IS CELL	20006
SUBROUTINE TRACE BEGINS AT CELL	21703

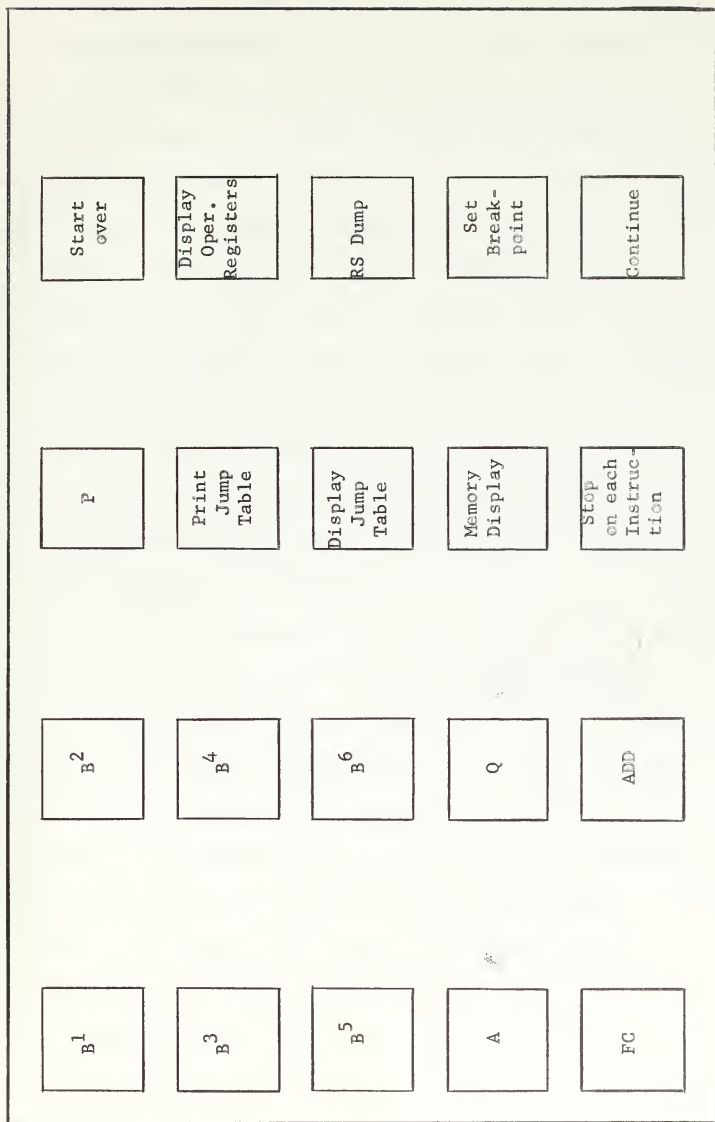


Figure 4. Keyboard #2

After displaying the introduction, Trace will wait for the programmer to take control at the DD 65 keyboard by pressing the CONTINUE button, which is shown in Figure 4. When the CONTINUE button is pressed, the object program will be allowed to run until the first jump instruction is reached. At that time, the contents of the operational registers will be displayed on the left tube (Figure 5), and the headings for the jump table will be displayed on the right tube (Figure 6). Each time the CONTINUE button is pressed, the object program will advance to the next jump instruction, and the operational register and jump table displays will be updated. The fourth column of the jump table will indicate the last four jump instructions executed in the object program. If an unexpected jump occurs, the table will show where the jump originated and the sequence that led up to it.

Keyboard #2 (Figure 4) provides the ability to change the contents of the operational registers, to examine the contents of memory, to record changes made in the program, to select and deselect various displays, to change the degree of control of program execution, and to restart the object program. The various functions are selected by pressing the desired button and deselected by pressing the button a second time. The keyboard light associated with the button indicates whether it is selected or deselected. Functions which might be conflicting are automatically deselected when a new function is selected. The functions which are initially selected

B1 00000 B2 00000 B3 00000
 B4 00000 B5 00000 B6 00000
 A 00000000 00000000
 Q 00000000 00000000
 P 000000
 PC 000 ADD 00000
 BREAK 00000

Figure 5. Operational Registers Display

TABLE OF EXECUTED JUMPS

FROM	TO	TIMES EXECUTED	LAST FOUR EXECUTED
20017	20030	00001	
20031	20014	00004	4
20018	20020	00011	LAST
20022	20014	00013	3
20016	20030	00005	2
20032	20020	00003	

Figure 6. Jump Table Display

when the program is started are PRINT JUMP TABLE, DISPLAY JUMP TABLE, DISPLAY OPERATIONAL REGISTERS, and CONTINUE (which is always active).

Control of the execution of the object program is exercised by the CONTINUE button as previously described. Variations are possible by the use of the STOP ON EACH INSTRUCTION and SET BREAKPOINT functions. With the STOP ON EACH INSTRUCTION function selected, the object program is stopped prior to the execution of each instruction and the displays updated. This provides the ability to step through portions of the program where closer examination is desired.

When the BREAKPOINT function is selected, the normal stops at each jump instruction are deleted; and when the CONTINUE button is pressed, the program runs continuously without stopping until the breakpoint address is reached or until the program is completed. The latter case will be discussed in a moment. Note that the program may be deliberately run to completion with no stops by selecting the BREAKPOINT function, but with the breakpoint address left zero.

The breakpoint address is set by first selecting the function and then typing the address on keyboard #1. The address will appear at the bottom of the operational register display as it is typed. If too many digits are typed, the address will be reset to zero and the breakpoint function deselected.

It should be noted that the octal number system is used for all displays and keyboard inputs.

The MEMORY DISPLAY function provides a means of inspecting constants or other desired portions of the object program. The function displays a block of twenty-one memory cells and may be selected at any time. After it is selected, the address of the first cell is typed on keyboard #1. This address will appear above the operational register display on the left tube. As the last digit of the address is typed, the operational register display will be cleared and the contents of that cell and the twenty successive memory cells will be displayed in its place. When the MEMORY DISPLAY is deselected, the operational register display is restored.

The DISPLAY JUMP TABLE and DISPLAY OPERATIONAL REGISTERS buttons provide the option of deleting the respective displays. Some computing time is saved by skipping these display generation routines. However, the time is normally insignificant compared with the dead time waiting for keyboard inputs, since the displays are only generated prior to a stop. The PRINT JUMP TABLE button can be used to delete the printing of a hard copy of the jump table at the end of the program. If both jump table functions are deleted, the table is not compiled. In this case, a significant amount of time is saved during continuous runs with the breakpoint set.

The left two columns on keyboard #2 and the button marked P are the register change functions. In each case, the desired register is selected and then the new contents for the register are typed on keyboard #1. As the contents are typed, they will be displayed on the screen. However, the format for typing is different for some of the registers.

For the index registers, B1 through B6, the number is typed in a normal manner from left to right, except that zeros at the lower end of the number need not be typed. For example, to put the number 02300 in one of the index registers, either 02300 or 023 may be typed. The function is deselected by pressing the selected button a second time or by striking a sixth key on keyboard #1.

For the A and Q registers, the procedure is slightly different. Quite often, the numbers desired in these registers amount to only a few digits in the lower end. Therefore, the order is reversed so that numbers are typed from right to left and zeros in upper portion of the register need not be typed. For example, to put the number 00000000 00002300 in the A register, the function is selected and then the number 0032 is typed on keyboard #1. To deselect the function, the A button is pressed a second time, or a seventeenth key is struck on keyboard #1.

The use of the function code, FC, and execution address, ADD, change functions are quite similar to that of the index registers. For the program address, P, register,

however, the full five digit address (including zeros) must be typed, and then a U or an L typed at the end to indicate either the upper or lower half of the address. As the new program address is typed, it will appear on the operational register display, and as the U or the L is typed, the remainder of the operational register display will change to correspond with the new program address.

The function code, execution address, and program address change functions should be used very carefully since they can disrupt the execution of the object program. The program address function, however, can be very useful for repeating or skipping sections of the object program.

The START OVER function provides an easy method of restarting the object program from the beginning. It may be used when the program has run to completion or at any intermediate time. To prevent accidental restarts, a two step operation is required. First, the button is pressed, and then any key on keyboard #1 is struck.

RSDUMP is a sub-program available on the system library. It provides a hard copy output of the object program in an octal format with mnemonics. When the RSDUMP button is pressed, the associated keyboard light will come on and remain on until the output is completed, and then the light will go out.

In the discussion of the keyboard functions, it should

be noted that the functions can only be used while execution of the object program is stopped, and that none of the functions except START OVER and P, will advance the object program until the CONTINUE button is pressed.

When the object program is run to completion, the word END will appear in the center of the left tube. All of the keyboard functions will be available at this time. However, the ones most likely to be used are MEMORY DISPLAY, to examine final values in the program; RSDUMP, to make a hard copy of the object program; or START OVER, to begin again. When the programmer is finished, the CONTINUE button is depressed to give control back to Monitor.

8. Trace Construction.

Figure 7 shows a simplified flow diagram of Trace with some functions omitted and with conditions as set initially in the program. The two loops in the program are clearly evident; one loop for jump instructions, and a second loop for the other instructions.

The program is entered initially from the calling statement in the object program. Trace then proceeds by pulling successive instructions out of the object program, one at a time. The type of instruction is determined by comparison with a table, and the instruction is then stored in the upper half (no matter whether it was an upper or lower half instruction) of a reserved cell

OBJECT PROGRAM

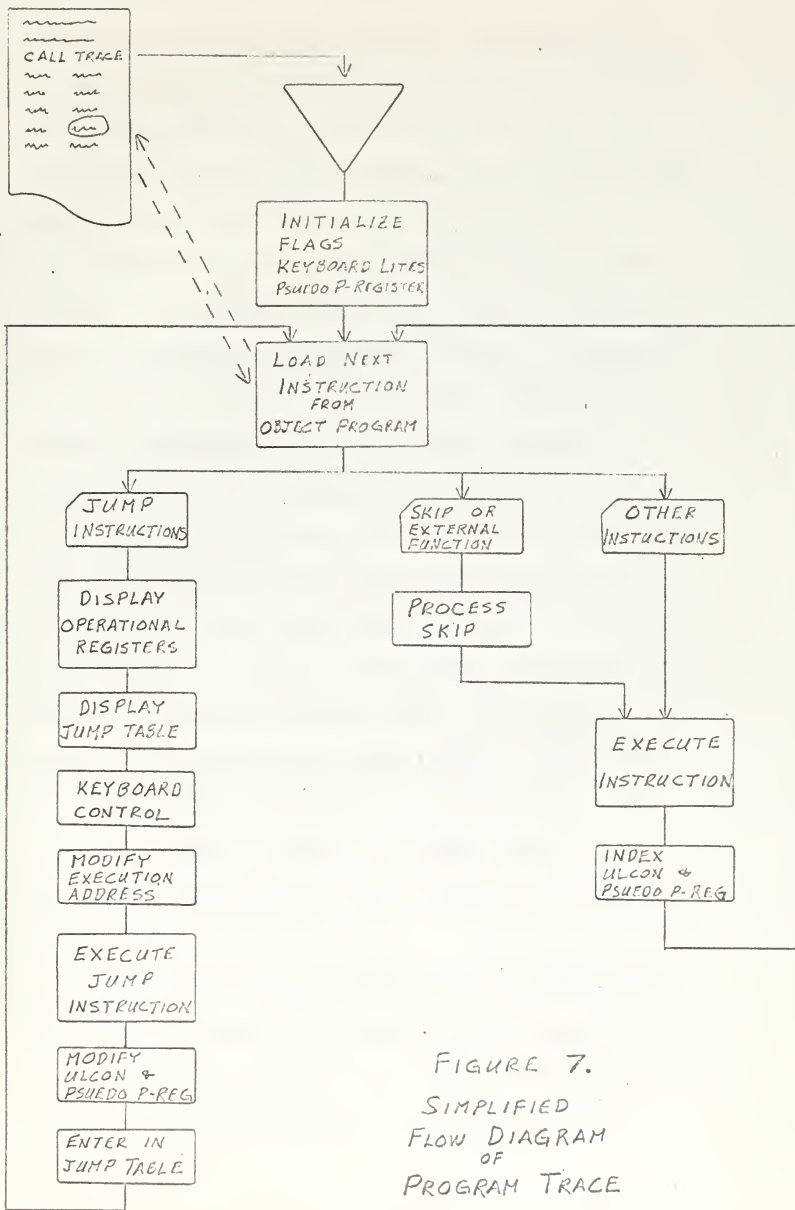


FIGURE 7.
SIMPLIFIED
FLOW DIAGRAM
OF
PROGRAM TRACE

either in the jump loop or the other instructions loop to await execution. An exception is that lower half skip and external function instructions are retained in the lower half, and some slight changes, indicated by the box marked Process Skip, are made to accommodate them.

Instructions other than jump instructions are then executed, and the upper-lower counter (ULCON), which keeps track of whether the instructions are upper or lower half instructions, and the psuedo P-register (PREG) are incremented as necessary. The loop is then completed by returning to load the next instruction from the object program.

For jump instructions, the operational register and jump table display routines and keyboard control are included in the loop. Also, before the jump instruction can be executed, its execution address must be modified so that control is retained in Trace. After the instruction is executed, the upper-lower counter (ULCON) and the psuedo P-register (PREG) are modified to correspond with the execution address and the type of jump (normal or return). The jump is then tabulated in the jump table and the loop is completed.

Although not shown in Figure 7, the display routines and keyboard control are bypassed when the breakpoint is set or when the DD 65 is not available. Also, the displays may be selectively bypassed by use of Keyboard Control.

In Figure 8, the flow diagram is expanded to show some additional functions and variations. For example, the display routines and Keyboard Control can be entered from the other instruction loop when the STOP ON EACH INSTRUCTION function is selected or when a breakpoint is reached. Also, the execution address of each jump is examined after it is tabulated in the jump table, and a jump to Monitor is used as the indication that the program is completed. If the program is completed, the jump table is printed, the word END is displayed, and Keyboard Control is entered.

Flow diagrams for Keyboard Control are shown in Figures 9 and 10. Keyboard Control is entered from one of three points marked by a, b, and c in Figure 8. Exit is by way of the CONTINUE button shown in Figure 9 and is always to the point from which the routine was entered. Keyboards #1 and #2 are alternately sensed until one or the other is hit. Figure 9 is the keyboard #2 portion of Keyboard Control, and Figure 10 is the keyboard #1 portion.

When keyboard #2 is hit, the input is decoded to determine which button was hit, and the function is selected or deselected by shifting a flag which is alternately positive and negative. If the function is selected, a return jump is made to the Master Deselect routine, which clears the flags and keyboard lights of conflicting functions. Those functions which have no conflicts do not use the Master Deselect routine.

OBJECT PROGRAM

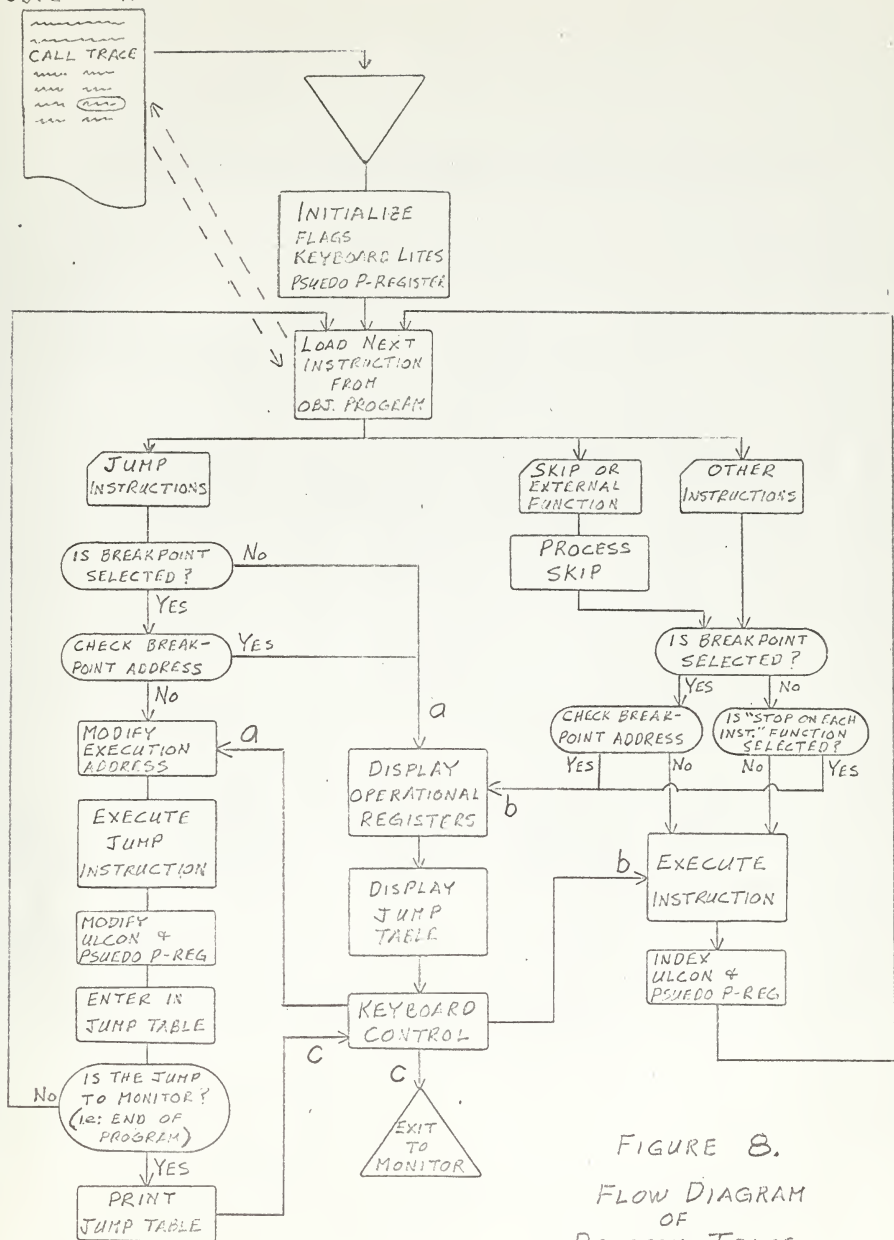


FIGURE 8.
FLOW DIAGRAM
OF
PROGRAM TRACE

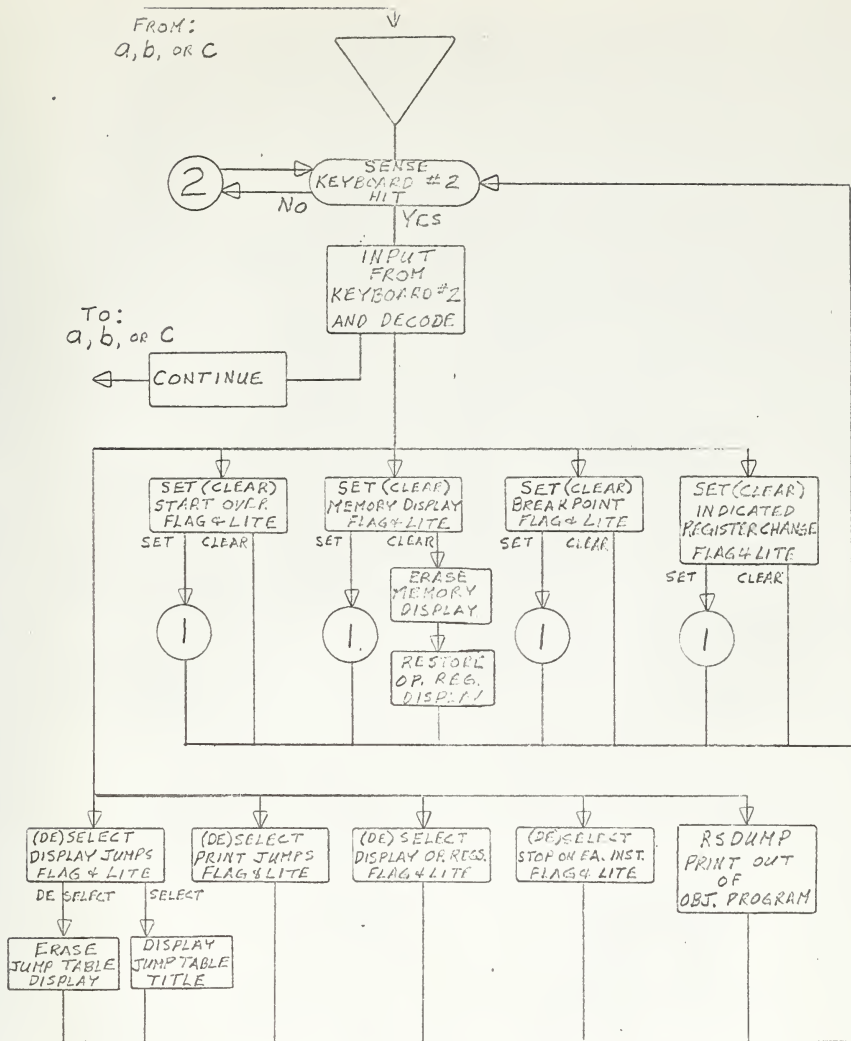
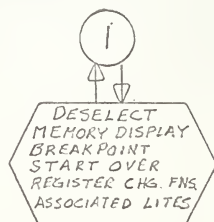


FIGURE 9.
KEYBOARD #2
PORTION OF KEYBOARD CONTROL



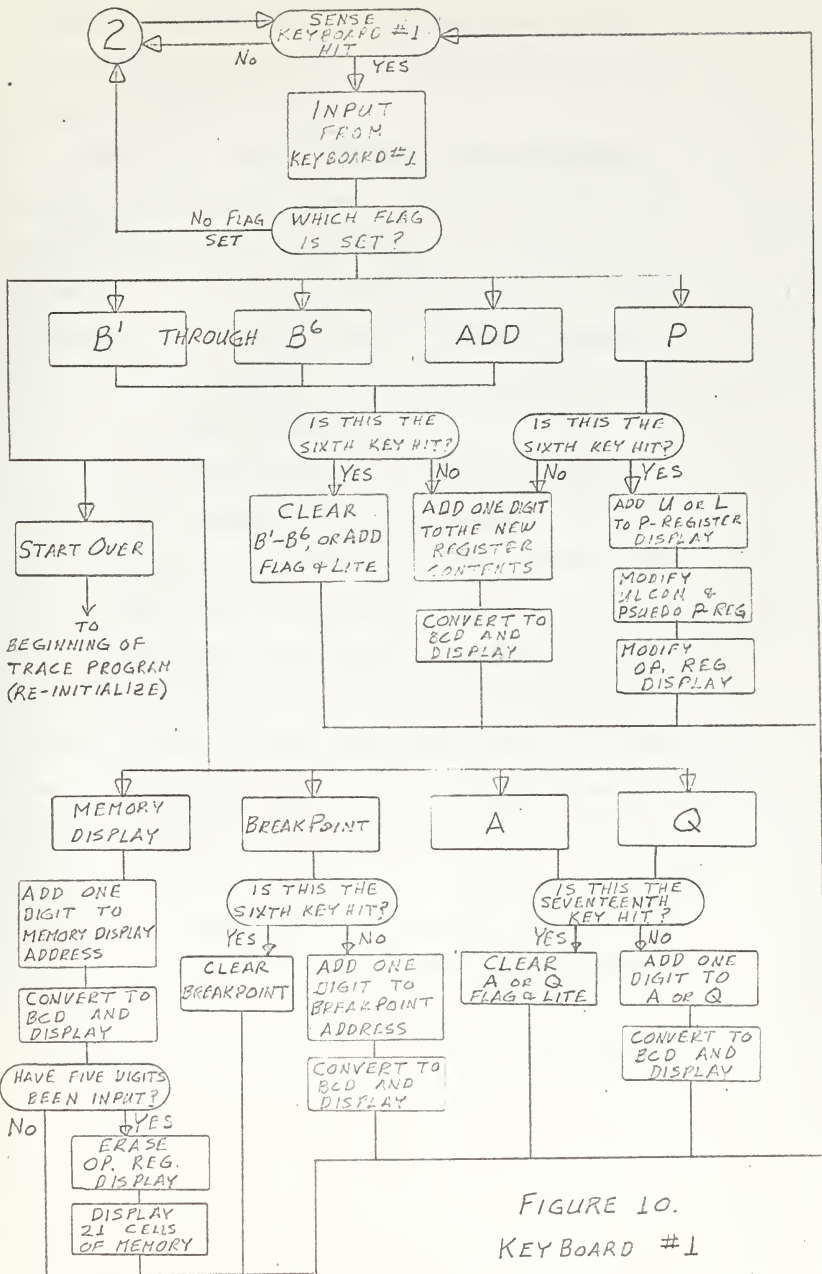


FIGURE 10.
KEYBOARD #1

PORTION OF KEYBOARD CONTROL

As shown in the diagram for the Memory Display and Jump Table Display functions, there are some additional display generation and erase features.

The register change functions, the Memory Display function, the Breakpoint function, and the Start Over function require inputs from keyboard #1. The keyboard #1 hits, which are in the form of a six bit BCD code, are input one at a time, and the control flags are checked to determine for which function they are intended. If no control flags are set, the input is discarded.

With two exceptions, the only inputs used are the octal digits 0 through 7. The first exception is that the Start Over function only checks that there has been a keyboard #1 hit. The second exception is the U or L needed by the Program Address change function. Otherwise, if any of the upper three bits of the BCD input are not zero, the input is treated as a zero. The reason being that the BCD codes for the other digits are just the digits themselves. The processing of the inputs by the various functions is shown in Figure 10. As each digit is received by the selected function, the desired number is assembled, both in binary form for use in the program and in BCD format for display. A count is kept of the number of digits input in order to determine when to deselect or execute the selected function.

Two routines are used for this purpose. A "48 bit registers" routine handles the 16 digit inputs for the A and Q registers, and a "15 bit registers" routine processes the five digit inputs for the other functions. When processing of a digit is completed, a return is made to the keyboard sensing wait loop to await the next keyboard hit.

Flow diagrams for other sections of Trace are included in Appendix II.

9. Conclusions.

The primary limitation of Program Trace is time, in terms of both running time and dead time. To calculate the running time for the program, consider that the breakpoint is set, the program is running, and the breakpoint has not yet been reached. Each non-jump, non-skip instruction in the object program requires the execution of from 43 to 45 instructions, and each skip or external function instruction requires the execution of from 43 to 48 instructions by Trace. For conditional jump instructions which are not satisfied, the number of instructions executed is 41. Executed jumps which are not a part of the object program require 59 instructions. These are jumps executed in library subprograms. They are executed under Trace control, but are not tabulated in the jump table.

For jump instructions in the object program, 86 to 101 instructions are executed by Trace. This includes the jump table tabulations and the end-of-program check, but it does not include display generation.

Thus, for example, an object program consisting of one-fifth jump instructions would have a running time of approximately 55 times normal when it is under Trace control with no display generation.

The display generation routines which are only executed before a stop require the execution of several hundred instructions. However, this time is negligible compared with the dead time waiting for keyboard inputs. Likewise, some of the Keyboard Control functions are relatively lengthy but their running time is small compared with the dead time. As long as the dead time is not utilized for other purposes, the length of the display routines and Keyboard Control is not important.

Although Program Trace was specifically designed for a non time-sharing system, it is constructed so that it could be incorporated into a simple time-sharing arrangement. For example, an area of computer memory could be reserved for Program Trace. The keyboard sensing wait loop, in which first one and then the other keyboard is checked until there is a hit, could be replaced by interrupts. Then, during the normal dead time, the computer could continue with other

jobs until interrupted by a keyboard hit. Since, in this arrangement, computer running time would be more important, it might also be desirable to transfer the display generation and Keyboard Control routines to the CDC 160 in order to take some of the load off of the CDC 1604.

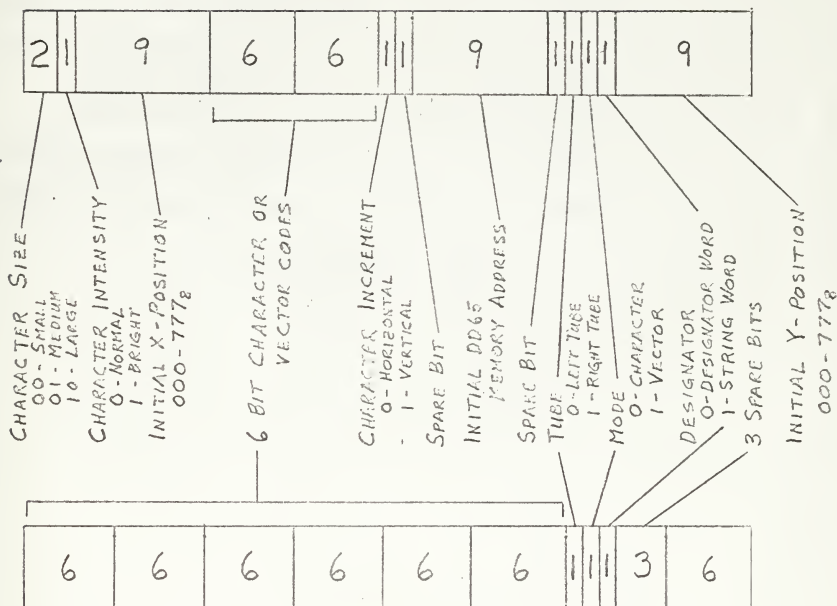
However, considering only the non time-sharing case, Program Trace can be very useful in many cases where the normal diagnostic methods become ineffective. The auto-monitoring of program execution can well be worth the investment of some additional computer time. Not only can an overall gain in efficiency of program diagnosis and checkout be realized, but in many cases, there will be no loss in computer time, since the need for repeated runs will be eliminated.

BIBLIOGRAPHY

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APPENDIX I

DD 65 DESIGNATOR WORD FORMAT



DD 65 STRING WORD FORMAT

DD 65 EXTERNAL FUNCTION CODES

Function	1604 Select Code
Select Keyboard 1 for Input	C7140
Keyboard 2 for Input	C7120
Select Track Ball "X"	C7102
Track Ball "Y"	C7104
Select Range Switch	C7110
Select Memory Update from 1604	C7010
Select Radar Target Data to 1604	C7002
Select Interruption:	
Keyboard 1 Hit	C7105
Keyboard 2 Hit	C7103
Radar Pulse Hit	C7141
Release Interrupt Request	C7111
Remove all Interrupt Selects	C7121

C = 1604 Data Channel Number

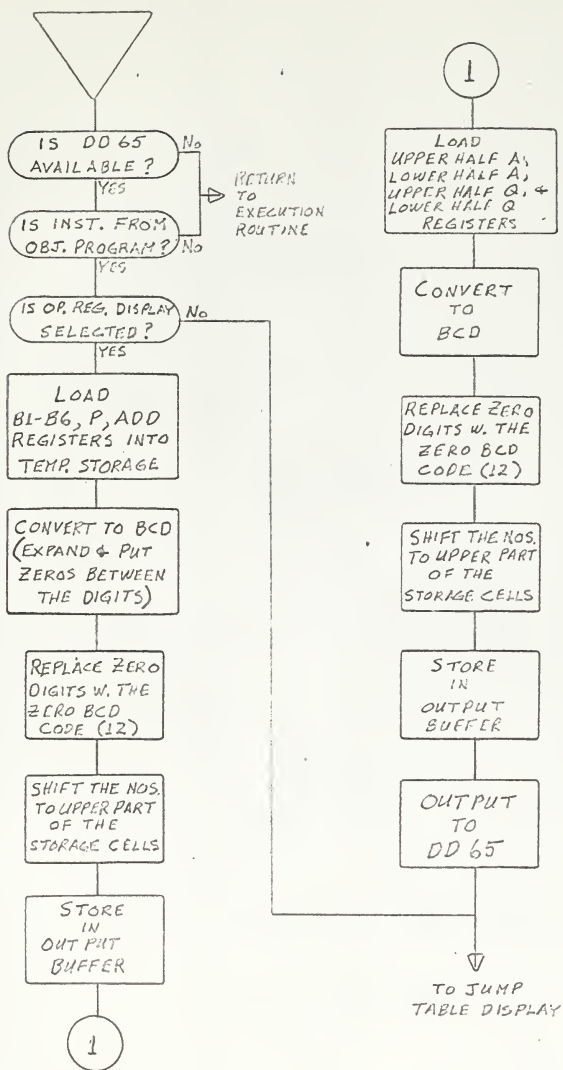
1604 SENSE CODES

Condition	Sense Code *
Keyboard 1 Hit	C7172
Keyboard 2 Hit	C7175
Keyboard 1 Not Hit	C7173
Keyboard 2 Not Hit	C7174
Tab Not Hit	C7157
Carriage Return Not Hit	C7166
Keyboard 1 Not Selected	C7167
Keyboard 2 Not Selected	C7037
DD 65 Interrupt	C7156
DD 65 from 1604 Selected	C7010
Radar to 1604 Selected	C7002

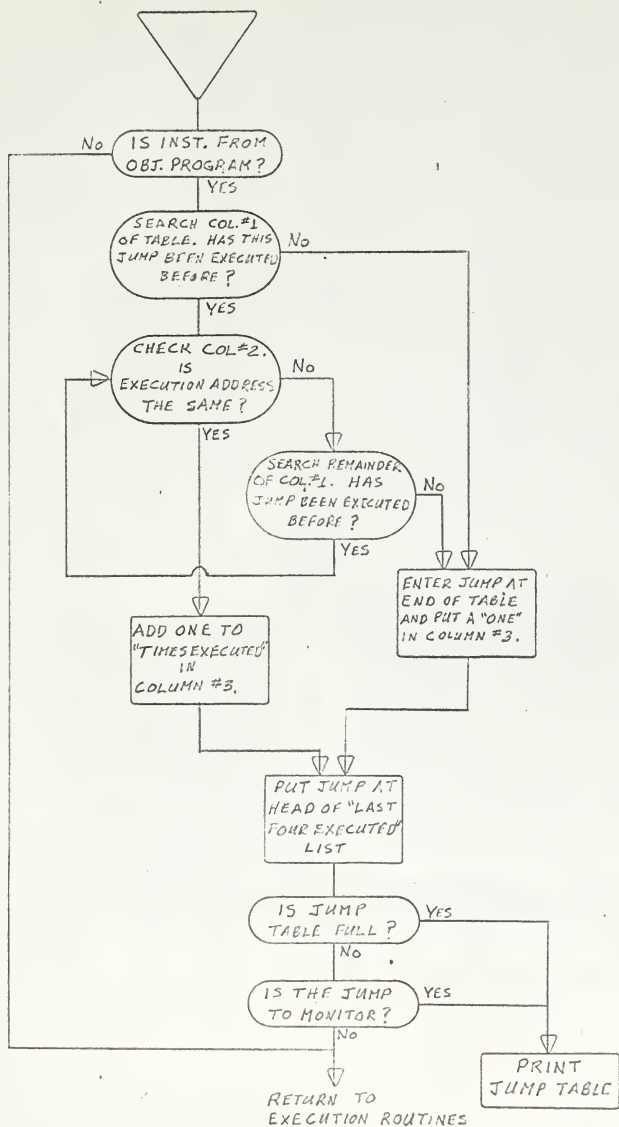
C = 1604 Data Channel Number

* Full Exit on a Positive Response.
 Half Exit on a Negative Response.

APPENDIX II



OPERATIONAL REGISTER DISPLAY



JUMP TABLE TABULATION

APPENDIX III PROGRAM TRACE

SUBROUTINE TRACE (ADDRESS)

```

DIMENSION BUF1(50), BUF2(50), BUF3(50)
CON=MSK1 = 000000007777777B; MSK2 = 7770000000000000B;
      MSK3 = 2000000000000000B; MSK4 = 7777777000000000B;
      MSK5 = 0000000000000000B; MSK6 = 0000000000070000B;
      MSK7 = 0000770000000000B; MSK8 = 000000007997999B;
      MSK9 = 7799000000000000B; MSK10 = 777000077777777B;
      MSK11 = 7777000000000000B; MSK12 = 0070000000000000B;
      MSK13 = 777700077777000B; MSK14 = 0000000007770000B;
      MSK15 = 0000000000000000B;
CON(IPASS = 5000000000000000B;
ZERO = 1212121212121212B;
UP = 000000000240000B;
END = 6525640000000000B;
ERASE= 0000000000001000B)

```

```

ULCON = 5252525252525252B;
LOW = 000000000004300000B;
BK = 121212121212000000B;
END1 = 57200000016321760B;

```

* OPERATIONAL REGISTERS OUTPUT BUF

```

CON(DDTB = 0012121212120000B;
BR1 = 5400620100001240B;
BR2 = 001212121200000B;
BR3 = 56062020041240B;
BR4 = 00121212120000B;
BR5 = 5400620400141140B;
BR6 = 56062050001140B;
AR = 540062060031140B;
ARR = 5400500000031740B;
ARL1 = 52121212120000B;
ARRR = 550121212120000B;
ARL = 54012121200361040B;
ARRR = 12121212120041740B;
ARR = 50012120041740B;
ARRR = 57401212120051740B;

```



```

6 PRI          54004700005416408)
  PR2          1212121222400008)
  PR3          54400000005616408)
  CON(PR3)    0012121200000618)
1 FC21        540066672006215408)
2 ADR1        00121212121200008)
3 ADR2        56206444006615408)
4 BK1         12121212120000008)
5 BK2         62516561420000008)
6 BK3         554000000007214408)

```

* INTRODUCTION

```

  CON(PREA)   00021212121200008)
1 PRE         23006363434300008)
2 PRE         65677145220000618)
3 PRE         61446243650000628)
4 PRE         51614400475100658)
5 PRE         24510047514600658)
6 PRE         3440304600010508)
  CON(PREB)   00121212121200008)
1 PRE         674343002235100618)
2 PRE         66236351001600658)
3 PRE         6451652220000618)
4 PRE         44465130006100648)
5 PRE         3440446500020108)
6 PRE         001212121200008)
  CON(PREC)   5165222200000008)
1 PRE         00612300616400228)
2 PRE         00622351616300658)
3 PRE         62514632423710758)
4 PRE         3440224004017508)

```

* JUMP TABLE TITLE

```

  CON(TITLE)  64004124444740228)
1 TITLE      65276563242340658)
2 TITLE      6243650046663408)
3 TITLE      3650236100653408)
4 TITLE      006642451004008)
5 TITLE      00000045612240258)
6 TITLE      4465229000004008)
7 TITLE      30502371011653208)
  CON(TI)    0000000023464008)
1 CON(TI)    46440000000004008)

```



```

3  TI = 35606551012653108,
4  TI = 63242385640040008,
5  TI = 0000000652740658,
6  TI = 6563242365644008,
   TI = 30406527013453108}

```

* JUMP TABLE OUTPUT BUF

```

CON(JT1) = 00121212121240008,
1  JT1 = 3550000014655708,
2  JT2 = 001212121240008,
3  JT3 = 3710000015252708,
4  JT4 = 001212121240008,
5  JT5 = 3050000015652708,
6  JT6 = 32100400014452708,
7  JT7 = 322300300015052708,
   CON(JT8) = 3223003000001460008,
1  JT9 = 32230030000015252708,
2  JT10 = 32704561015252708}

```

* MEMORY DUMP OUTPUT BUF

```

CON(DP1) = 12121212120000008,
1  DP1 = 64244447000000008,
2  DP2 = 5620000010013208,
3  DP3 = 6164451652202208,
4  DP4 = 5400000131012608,
5  DP5 = 46453565432302208,
9  DP6 = 5000006313142608,
7  DP7 = 121212120000008,
   CON(DP8) = 54200000022012008,
1  DP9 = 12121212121200008,
2  DP10 = 51401212022412008,
3  DP11 = 1212121212120008,
4  DP12 = 57001212023012008}

```

* INITIALIZE

```

1AA  SIU1(DEX)
      STA(TEMA)
      LDA(ULCON)
      STA(BREAK)
      LAC(ULCON)
      STA(DTSJUM)
      STA(B1FLG)
      STA(B2FLG)
      STA(B3FLG)
      STA(B5FLG)

```

```

      SIL2(DEX)
      SIQ(TEMQ)
      STA(OPREG)
      STA(STPRINT)
      STA(DUMP)
      STA(PRTJUM)
      STA(B2FLG)
      STA(B4FLG)
      STA(B6FLG)

```

```

      STORE INDEX 1,2
      STORE A AND Q
      INITIALIZE
      CONTROL FLAGS

```

.....


```

STA (AFLG)
STA (FCFLG)
STA (PFLG)
STA (ELDW)
SXXF (77508B)
SXXF (77527B)
SXXF (77441B)
SXXF (77421B)
SXXF (77408B)
SXXF (77404B)
SXXF (77308B)
SXXF (77308B)
SXXF (77308B)
SXXF (77318B)
SXXF (77321B)
SXXF (77321B)
SXXF (77241B)
SXXF (77341B)
SXXF (77408B)
SXXF (77508B)
SXXF (77518B)
SXXF (77540B)
SXXF (77708B)
STA (PREG)
LDA (ULCON)
SLJ4 (1E)

```

```

STA (QFLG)
STA (ADFLG)
STA (START)
ENI (0)
ISKI (10)
ISKI (10)
ISKI (10)
ISKI (10)
ISKI (10)
ISKI (10)
ISKI (10)
ISKI (10)
ISKI (10)
ISKI (10)
ISKI (10)
ISKI (10)
LDA (ADDRESS)
ENI2 (1)
STA (UPLOW)
ZRO (0)

```

INITIALIZE
LIGHTS

INITIALIZE PSUEDO P-REGISTER
INITIALIZE UPPER-LOWER COUNTER

* CONTROL

```

1A LDA7 (PREG)
   LSH (UPLOW)
   SCL (MSKT)
   SCL (6)
   LIL (TM)
   SAL (L+2)
   LDA (UPLOW)
   SIL1 (ADDRESS)

```

```

LOAD NEXT INST FROM PROGRAM
STORE IN UPPER HALF OF TM
STORE FUNCTION CODE IN TM
LOOK IN TABLE
SAVE ADD, GO TO EXEC ROUTINE

```

* TABLE

```

1L ZRO (0)
   ZRO (0)
   ZRO (0)
   ZRO (0)

```

```

SLJ (10)
SLJ (10)
SLJ (10)
SLJ (10)

```

```

ZRS
ORS
LRS
ENS

```


ZRO(0)
ZRO(0)
ZRO(0)
ZRO(0)
ZRO(0)
ZRO(0)
ZRO(0)
ZRO(0)
ZRO(0)
ZRO(0)

SLJ(1S)
SLJ(1S)
SLJ(1S)
SLJ(10)
SLJ(10)
SLJ(10)
SLJ(10)
SLJ(1S)
SLJ(1S)
SLJ(1J)
SLJ(10)

EQS
MTH
RAD
RSB
RAO
RSO
EXF
SLJ
SEV

* JUMP INSTRUCTIONS

1J	LDQ(MSK2) LQC(MSK2) STA(TEXEC) LDA(BREAK) LDA(BRKPT) MEQ(PREG) SLJ4(1D) SLDQ(TEM0) SLJ1(STORE) SLJ1(DEX) SLJ(L+1) ENI(0) SLJ(L+1) SLJ(0) SLJ(DEX) SLJ(0) SLJ(DEX) LDA(MSK3) MEQ(TEM) LDA(PREG) SAU7(ADDRESS) AP2(L+1) STA(UPLOW) LDA(UPLOW) AP3(L+1) STA(UPLOW) LDA(PREG) LDA(ADDRESS) SSK(PRTJUM) SSK(DISJUM) SLJ(1A) SLJ(DEX) SLJ(STORE) LSSK(UPLOW) RAD(PREG)	LDL(TEM) ADL(TEXEC) ENI(0) AP3(L+3) LDQ(MSK5) SLJ(L+2) ENI(0) LDA(TEMA) SLJ2(STORE) SLJ2(DEX) SLJ(4J) ENI(0) SLJ(DEX) LIL2(STORE) LIL2(MSK3) LDQ(MSK3) SLJ(2J) INA(1) LDA(UPLOW) ALS(1) SLJ(3J) ENI(0) ALS(1) ENI(0) STA(TEMP) STA(PREG) SLJ(1T) SLJ(1T) ZRO(0) SLJ2(DEX) LIL2(STORE) SLJ(1A) SLJ(1A)	LOAD JUMP INST MOD ADDRESS IS BREAKPOINT SET CHECK BREAKPOINT GO TO DISPLAY LOAD A AND Q SAVE TRACE INDEX 1,2 LOAD PROGRAM INDEX 1,2 EXEC INST NO JUMP GO TO 4J SAVE INDEX 1,2 IF NOT RET JUMP GO TO 2J PROCESS RET JUMP PROCESS JUMP PROCESS JUMP SAVE LAST INST ADDRESS CHANGE P-REGISTER GO TABLE JUMPS GO TABLE JUMPS PROCESS NON-JUMP INCREASE P-REGISTER
2J			
3J			
4J			

* SKIP INSTRUCTIONS

15 SSK(UPLOW)
 LDA7(PREG)
 ADD(PASS)
 SLJ(30)
 SLJ(10)
 SCL(MSK4)
 STA(2EXEC)
 ZRO(0)

IF UPPER INST GO TO 10
 IF LOWER INST
 MOD 2EXEC

* OTHER INSTRUCTIONS

20 ZRO(0)
 LQC(MSK1)
 LDO(MSK1)
 STA(2EXEC)
 LDA(BREAK)
 SSK(STPINT)

MOD 2EXEC
 IS BREAKPOINT SET
 IS STOP ON EACH INST. FLAG SET

30 SLJ(L+4)
 LDA(BRKPT)
 MEQ(PREG)
 SLJ4(1D)
 SIU1(STORE)
 LDO(TEMQ)
 ZRO(0)

CHECK BREAKPOINT
 GO TO DISPLAY
 SAVE TRACE INDEX 1,2
 LOAD PRG INDEX 1,2

2EXEC ZRO(0)
 STA(TEMA)
 ENI(0)
 SIU1(DEX)
 LIU1(STORE)
 AJP3(L+1)
 STA(UPLOW)
 SIU1(DEX)
 LIU1(STORE)
 SSK(UPLOW)
 RAD(PREG)

EXEC INST. NO SKIP GO TO 40
 PROCESS SKIP

40 STA(UPLOW)
 SIU1(DEX)
 LIU1(STORE)
 AJP3(L+1)
 STA(UPLOW)
 SIU1(DEX)
 LIU1(STORE)
 SSK(UPLOW)
 RAD(PREG)

PROCESS NON SKIP

50 STA(UPLOW)
 RAD(PREG)

IF LOWER INST
 INCR P-REGISTER

* DISPLAY OPERATIONAL REGISTERS

10 SLJ(N)
 SSK(5E)
 ENA(L)
 AJP3(1D)
 SUB(OPREG)
 SSK(STORE)
 STA(HOLD1)
 STA(HOLD2)

EXIT/ENTRY
 SKIP DISPLAY IF DD65 NOT AVAIL
 IF ABOVE OR BELOW
 MAIN PROGRAM
 SKIP DISPLAY
 SKIP IF NOT SELECTED
 PREPARE CONTENTS
 OF INDEX REGISTERS,
 P REGISTER AND

20	STA(HOLD1) SIL4(HOLD4) SIL6(HOLD6) STA(HOLD7) STA(HOLD8) LDQ(MSK6) ENI113 ALS(13) ADL2(HOLD) LDQ(MSK7) ENI114 ENA(O) MEQ(HOLD) LDL(ZERO) LJPI(30) ALS112 ALS2(8) ENI2(1) LDA2(HOLD) INI1(2) ISK2(6) LDA(TEMA) ALS(24) LDQ(TEQ) OLS(24) ENI2(7) LDQ(MSK8) ENI116 ALS(13) ADL2(HOLD) LDQ(MSK9) ENI117 SLJ(L+2) ENA(O) MEQ(HOLD) LDL(ZERO) LJPI(60) ENI1(O) ISK2(4) ENI2(1) LDA2(HOLD) ALS(12) INS1(2) INSK2(1) LDQ(MSK11) ENI2(1) LDQ(MSK12)	SIL3(HOLD3) SIL5(HCLD5) LDA(PREG) LDA(ADRES) ENI2(1) LDL2(HOLD) ENI(O) QRS(3) IJP1(L-1) STA(HOLD) ENI(O) QRS(6) SLJ(L+2) RAD(HOLD) LDA(HOLD) STA2(HOLD) SLJ(20) ENI(O) STA(DCTB) ENI(O) SLJ(40) STA(HOLD2) STA(HOLD1) STO(HOLD4) STO(HOLD3) EXF(770108) LDL2(HOLD) ENI(O) QRS(3) IJP1(L-1) STA(HOLD) ENA(O) ZRO(O) OLS(6) SLJ(L+2) RAD(HOLD) LDA(HOLD) STA2(HOLD) SLJ(50) ENI(O) SCL(MSK11) STA1(ARL1) ENI(O) SLJ(70) ENI1(1) LDL2(HOLD) LDQ(MSK10)	EXECUTION ADDRESS FOR DISPLAY
30		PUT REGISTER CONTENTS IN BCD FORMAT	
40		REPLACE ZEROS WITH THEIR EQUIVALENT BCD CODE NO ZERO	
50		STORE IN OUTPUT BUF	
60		PREPARE CONTENTS OF A AND Q REGISTERS FOR DISPLAY	
70		PUT REGISTER CONTENTS IN BCD FORMAT	
80		REPLACE ZEROS WITH THEIR EQUIVALENT BCD CODE NO ZERO	
		STORE	
		IN	
		OUTPUT	

ADL1(ARL1)
 INI1(2)
 ISK2(4)
 LDA(HOLD7)
 STA(PR2)
 JP2(L+2)
 RAD(PP2)
 LDA(HOLD8)
 LDQ(MSK12)
 ARS(3)
 ADL(TEM)
 QLS(3)
 ARS(9)
 INA(6TB)
 SCL(MSK2)
 LDQ(MSK7)
 QRS(6)
 MEQ(FC1)
 LDL(ZERO)
 QLS(6)
 IJP1(L-3)
 OUT1(DOTB)

9D

* DISPLAY JUMP TABLE

SSK(QJSJUM)
 SLJ(1K)
 LDA(L1)
 SAU(100B)
 ENI3(3)
 LDQ(MSK6)
 QRS(3)
 ADL1(BUF1)
 STA(JT1)
 ENI3(4)
 ENA(0)
 MEQ(JT1)
 LDL(ZERO)
 IJP3(L-3)
 LDQ(MSK6)
 QRS(3)
 ADL1(BUF2)
 STA(JT2)
 ENI3(4)
 ENA(0)
 MEQ(JT2)
 LDL(ZERO)

48

STAI(ARL1)
 ENI(0)
 SLJ(8D)
 ALS(6)
 LDA(UL)
 LDA(LWR)
 SLJ(L+2)
 RAD(PR2)
 STA(ADR1)
 LDL(TEM)
 QLS(3)
 ARS(3)
 ADL(TEM)
 INA(6TB)
 STA(PL)
 ENI(2)
 ENA(0)
 SLJ(L+2)
 RAD(FC1)
 ENA(0)
 ENI1(32)
 LIL2(STORE)

.....

CHECK FOR
 LOWER OR
 UPPER INST
 STORE P REG
 PREPARE FUNCTION CODE
 FOR DISPLAY

.....

.....

SLJ(L+2)
 ZRO(0)
 LIL1(L1)
 SIL3(HCLD)
 ENI(0)
 LDL1(BUF1)
 ALS(3)
 IJP3(L-1)
 LDQ(MSK7)
 ENI(0)
 QRS(6)
 SLJ(L+2)
 RAD(JT1)
 ENI3(3)
 LDL1(BUF2)
 ALS(3)
 IJP3(L-1)
 LDQ(MSK7)
 ENI(0)
 QRS(6)
 SLJ(L+2)
 RAD(JT2)

.....

.....

IJP3(L-3)
LDQ(MSK6)
QRS(3)
ADL1(BUF3)
STA(JT4)
ENA(0)
MEO(JT4)
LDL(ZERO)
IJP3(L-3)
STA(JT)
ALS(12)
STA(JT2)
ALS(JT2)
STA(JT4)
STA(HOLD1)
ENQ(7778)
SUB(HOLD1)
STA(HOLD1)
SCL(HOLD2)
ADL(HOLD1)
LDA(JT3)
SCL(MSK15)
STA(JT3)
SCL(HOLD2)
ADL(HOLD1)
ENA(2)
LDA(HOLD1)
ENQ(0)
DVI(HOLD8)
ENA(13068)
STA(HOLD1)
RAD(JT1)
RAD(HOLD1)
RAD(JT3)
RAD(HOLD1)
RAD(JT5)
REF(777108)
TSKI(N)
ENA(1278)
ALS(3)
SCL(MSK13)
STA(JT6)
SUB(L3)
SCL(MSK13)
STA(JT7)

100B

ENI3(3)
LDL1(BUF3)
ALS(3)
IJP3(L-1)
LDQ(MSK7)
ENI(0)
QRS(6)
SLJ(L+2)
RAD(JT4)
LDA(JT)
LNA(40008)
LDA(JT2)
LNA(40008)
LDA(JT4)
LNA(40008)
ENA(0)
ENA(1278)
STO(HOLD2)
ALS(3)
LDA(JT1)
SCL(MSK15)
STA(JT1)
SCL(HOLD2)
ADL(HOLD1)
LDA(JT5)
SCL(MSK15)
STA(JT5)
STA(HOLD8)
LNA(4508)
STA(HOLD1)
RAD(HOLD1)
SUB(HOLD1)
ALS(12)
ENA(4)
ALS(12)
ENA(4)
ALS(12)
ENI3(6)
OUT3(JT)
SLJ(10DA)
SUB(L4)
LQC(HOLD2)
ADL(JT6)
ENA(1278)
ALS(3)
ADL(JT7)
ENA(1278)

PUT TIMES EXECUTED
IN BCD FORMAT

PREPARE TAGS OF LAST
FOUR JUMPS FOR OUTPUT



SUB(L2)
 SCL(MSK13)
 STA(JT8)
 SUB(L1)
 SCL(MSK13)
 STA(JT10)
 OUT(JT10)
 OUT(JT6)
 SLJ(K)

* MEMORY DISPLAY

11D
 XF(177010B)
 ENA(1000B)
 ENI(52B)
 OUT(ERASE)
 RAD(ERASE)
 ENI(200B)
 ENA(1320B)
 STA(HOLD4)
 ENI(1230B)
 STA(HOLD6)
 ENI(17B)
 LDA(DPT)

12D
 LDA(N)
 ARS(24)
 ENI(20)
 LDO(MSK8)
 ENI(6)
 ENA(53)
 ADL2(HOLD)
 LDO(MSK9)
 ENI(7)
 ENA(L+2)
 ENA(O)

14D
 MEQ(HOLD)
 LDO(ZERO)
 LDO(JP(140))
 SCL(MSK11)
 SCL(AZ(DP9))
 LDO(MSK13)
 ADD(HOLD1)
 STA2(3)
 ENI(3)

ALS(3)
 ADL(JT8)
 ENA(127B)
 ALS(3)
 ADL(JT10)
 ENI(5)
 L13(HOLD)
 ZRO(O)

ENI(O)
 STA(ERASE)
 ENI(O)
 ENA(20000B)
 IJP(LP3)
 OUT(HOLD1)
 ALS(12)
 ENA(1324B)
 STA(HOLD5)
 ALS(12)
 SHL3(HOLD7)
 ENI(O)
 STA(L+1)
 STA(HOLD2)
 ENI(O)
 LDO2(HOLD)
 ENI(O)
 QMS(3)

IJP(L-1)
 STA(HOLD)
 ENA(O)
 ZRO(O)
 OLS(6)
 SLJ(L+2)
 RAD(HOLD)
 LDO(HOLD)
 ALS(12)
 LDO(MSK11)
 LDO(L+2)
 ARS(12)
 ADL2(DP10)
 SCL(MSK15)
 INI2(1)
 SLJ(13D)
 LDO(MSK6)

ERASE OPERATIONAL REGISTER DISPLAY
 DISPLAY TABLE HEADINGS

PREPARE CONTENTS OF DUMP ADDRESS FOR DISPLAY

PUT IN BCD FORMAT

REPLACE ZEROS WITH THEIR EQUIVALENT BCD CODE

STORE IN OUTPUT BUFFER

PREPARE DUMP
ADDRESS FOR
DISPLAY

LD(LDPT)
 ALS(13)
 ADL(DPT)
 STA(HOLD3)
 ENA(14)
 ENA(10)
 MEQ(HOLD3)
 LD(LZER0)
 IJ(LZER3)
 ALS(12)
 LDQ(MSK13)
 SCL(MSK15)
 ADD(HOLD4)
 LDA(HOLD5)
 LDA(HOLD6)
 ENA(16)
 ENA(14B)
 STA(HOLD8)
 LDA(HOLD8)
 LDA(HOLD8)
 LDA(HOLD1)
 SCL(MSK13)
 STA(HOLD1)
 IJP(12D)
 LIL3(HOLD7)

* KEYBOARD 2

1K
 EXE(77175B)
 EXE(77120B)
 ENQ(40B)
 AJP(11K)
 LD(LKYBD2)
 AJP(LL+1)
 ADL(KYBD2)
 STA(HOLD)
 MEQ(KYBD2)
 LDA(HOLD)
 SAU(LL+1)
 LIU(LN)
 LDA(HOLD)
 SAU(LL+1)
 LIU(LN)
 LIU(LN)
 SIU(LL+1)
 SLJ(IN)

2K
 * TABLE

EN(9)
 QRS(13)
 IJP(LL-1)
 LDQ(MSK7)
 EN(0)
 QRS(6)
 SLJ(LL+2)
 RAD(HOLD3)
 LDA(HOLD3)
 STA(DP7)
 STA(DP8)
 LD(LDP8)
 ADD(HOLD1)
 STA(DP8)
 RAD(DP10)
 RAD(DP12)
 OUT(DP7)
 ALS(12)
 RAD(HOLD4)
 RAD(HOLD5)
 RAD(HOLD6)
 INA(-20B)
 SCL(MSK4)
 RAD(DPT)
 EN(0)
 LIL2(STORE)

OUTPUT TO DD65

SENSE KYBD 2 HIT
 INPUT KYBD 2 HIT
 PROCESS KYBD 2 HIT

GO TO SELECTED CONTROL FN

SLJ(173K)
 SLJ(174K)
 SLJ(175K)
 SLJ(176K)
 SLJ(177K)
 SLJ(178K)
 SLJ(179K)
 SLJ(180K)
 SLJ(181K)
 SLJ(182K)
 SLJ(183K)
 SLJ(184K)
 SLJ(185K)
 SLJ(186K)
 SLJ(187K)
 SLJ(188K)
 SLJ(189K)
 SLJ(190K)
 SLJ(191K)
 SLJ(192K)
 SLJ(193K)
 SLJ(194K)
 SLJ(195K)
 SLJ(196K)
 SLJ(197K)
 SLJ(198K)
 SLJ(199K)
 SLJ(200K)
 SLJ(201K)
 SLJ(202K)
 SLJ(203K)
 SLJ(204K)
 SLJ(205K)
 SLJ(206K)
 SLJ(207K)
 SLJ(208K)
 SLJ(209K)
 SLJ(210K)
 SLJ(211K)
 SLJ(212K)
 SLJ(213K)
 SLJ(214K)
 SLJ(215K)
 SLJ(216K)
 SLJ(217K)
 SLJ(218K)
 SLJ(219K)
 SLJ(220K)
 SLJ(221K)
 SLJ(222K)
 SLJ(223K)
 SLJ(224K)
 SLJ(225K)
 SLJ(226K)
 SLJ(227K)
 SLJ(228K)
 SLJ(229K)
 SLJ(230K)
 SLJ(231K)
 SLJ(232K)
 SLJ(233K)
 SLJ(234K)
 SLJ(235K)
 SLJ(236K)
 SLJ(237K)
 SLJ(238K)
 SLJ(239K)
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 SLJ(250K)
 SLJ(251K)
 SLJ(252K)
 SLJ(253K)
 SLJ(254K)
 SLJ(255K)
 SLJ(256K)
 SLJ(257K)
 SLJ(258K)
 SLJ(259K)
 SLJ(260K)
 SLJ(261K)
 SLJ(262K)
 SLJ(263K)
 SLJ(264K)
 SLJ(265K)
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 SLJ(267K)
 SLJ(268K)
 SLJ(269K)
 SLJ(270K)
 SLJ(271K)
 SLJ(272K)
 SLJ(273K)
 SLJ(274K)
 SLJ(275K)
 SLJ(276K)
 SLJ(277K)
 SLJ(278K)
 SLJ(279K)
 SLJ(280K)
 SLJ(281K)
 SLJ(282K)
 SLJ(283K)
 SLJ(284K)
 SLJ(285K)
 SLJ(286K)
 SLJ(287K)
 SLJ(288K)
 SLJ(289K)
 SLJ(290K)
 SLJ(291K)
 SLJ(292K)
 SLJ(293K)
 SLJ(294K)
 SLJ(295K)
 SLJ(296K)
 SLJ(297K)
 SLJ(298K)
 SLJ(299K)
 SLJ(300K)

* CONTROL FLAGS

3K SSH(OPREG)
 EXF(775058)
 EXF(775048)
 SLJ(17)
 SLJ(18)
 SSH(DI5JUM)
 ENI(01)
 ENI(0115)
 SAU(1008)
 SAU(1008)
 EXF(77410B)
 EXF(77411B)
 ENA(1061B)
 STA(EASE)
 OUT(EASE)
 RAD(EASE)
 SLJ(14)
 SSH(DUMP)
 SLJ4(23K)
 SLJ(1D+1)
 SLJ4(23K)
 SSH(DUMP)
 STA(DPT)
 ENA(4)
 EXF(77420B)
 SSH(PRTJUM)
 EXF(774048)
 EXF(774058)
 SSH(BREAK)
 SLJ4(23K)
 EXF(77520B)
 ENA(4)
 SLJ(1K)
 SLJ4(23K)

3K

4K
 4KAA

4KAB
 4KA

5K

6K

7K

SLJ(12K)
 SLJ(6K)
 SLJ(4K)
 SLJ(5K)
 SLJ(9K)
 SLJ(11K)
 SLJ(3K)
 SLJ(9K)
 SLJ(10K)
 SLJ(10K)

SLJ(L+2)
 SLJ(1K)
 SLJ4(23K)
 ZRO(0)
 SLJ(4KA)
 EXF(77010B)
 OUT(1TITLE)
 LDA(STORE)
 ENI(10)
 SLJ(10CA)
 ENI(1500B)
 ALS(9)
 EXF(77010B)
 ENA(2000B)
 JPL(L-1)
 ZRO(0)
 SLJ(L+3)
 ZRO(0)
 ZRO(0)
 ZRO(0)
 ENA(0)
 STA(DP)
 STA(BRKDEX)
 SLJ(1K)
 SLJ(L+2)
 SLJ(1K)
 SLJ(1K)
 SLJ(1K)
 SLJ(L+5)
 ZRO(0)
 SSH(BREAK)
 STA(BRKDEX)
 ZRO(0)
 ZRO(0)

B1 -- PRINT JUMPS
 B3 -- DISPLAY JUMPS
 B5 -- MEM DUMP
 A -- STOP ON EA INST
 B2 -- START OVER
 B4 -- DISPLAY OP REGS
 B6 -- DISPLAY FLOW DIA
 Q -- SET BREAK
 ADD -- CONT

DISPLAY OPERATIONAL REGS FLAG

DISPLAY JUMP TABLE FLAG

ERASE
 JUMP
 TABLE
 DISPLAY

MEMORY DUMP FLAG

PRINT JUMP TABLE FLAG

BREAKPOINT FLAG

8K	SLJ(1K) ENA(1AA) LDA(ADDRESS) CALL RSDUMP(NL,NAN,3) EXH(77511B) LJ4(23K) EXF(77440B) EXF(77441B) LJ(1D) LJ(START) EXF(77503B) LJ4(23K) SSH(START) LJ(1K) SSH(PFLG) EXF(77403B) LJ4(23K) SSH(PELG) STA(BRKDEX) LJ(1K) SSH(B1FLG) EXF(77203B) LJ4(23K) SSH(B1FLG) STA(BRKDEX) LJ(1K) SSH(B3FLG) EXF(77205B) LJ4(23K) SSH(B3FLG) STA(BRKDEX) LJ(1K) SSH(B5FLG) EXF(77211B) LJ4(23K) SSH(B5FLG) STA(BRKDEX) LJ(1K) SSH(AFLG) EXF(77221B) LJ4(23K) SSH(AFLG) ENA(O) SSH(FCFLG)	ZRO(O) EXF(77510B) STA(NAN) STA(NL) SLJ(1K) LJ(L+3) ZRO(O) SLJ(1K) SLJ(1K) ZRO(O) SLJ(L+2) SLJ(1K) ZRO(O) EXF(77502B) ZRO(O) SLJ(L+2) ZRO(O) ENA(4) EXF(77402B) ZRO(O) SLJ(L+2) SLJ(1K) ZRO(O) ENA(4) EXF(77202B) ZRO(O) SLJ(L+2) ZRO(O) ENA(4) EXF(77204B) ZRO(O) SLJ(L+2) ZRO(O) ENA(4) EXF(77210B) ZRO(O) SLJ(L+2) ZRO(O) ENA(4) EXF(77208) ZRO(O) SLJ(L+2) ZRO(O) ENA(4) EXF(77210B) ZRO(O) SLJ(L+2) ZRO(O) ENA(4) EXF(77208) ZRO(O) SLJ(L+2)	RSDUMP FLAG STOP ON EA INST FLAG CONTINUE START OVER FLAG P-REG FLAG B1 FLAG B3 FLAG B5 FLAG A-REG FLAG FUNCTION CODE FLAG
----	--	--	---

18K EXF(77241B)
 SLJ4(23K)
 SSH(FCFLG)
 STA(BRKDEX)
 SSH(B2FFLG)
 SLJ(1K)
 EXF(77303B)
 SLJ4(23K)
 SSH(B2FLG)
 TA(BRKDEX)
 SLJ(1K)
 SSH(B4FLG)
 EXF(77305B)
 SLJ4(23K)
 SSH(B4FLG)
 STA(BRKDEX)
 SLJ(1K)
 SSH(B6FLG)
 EXF(7731B)
 SLJ4(23K)
 SSH(B6FLG)
 STA(BRKDEX)
 SLJ(1K)
 SSH(QFLG)
 EXF(77321B)
 SLJ4(23K)
 SSH(QFLG)
 STA(BRKDEX)
 SLJ(1K)
 SSH(ADFLG)
 EXF(77341B)
 SLJ4(23K)
 SSH(ADFLG)
 STA(BRKDEX)
 SLJ(1K)

* MASTER DESELECT

23K SLJ(N)
 ENA(O)
 STA(BRKPT)
 LDA(BK)
 STA(BK1)
 ENI(13)
 EXF(77010B)
 OUT1(BK1)
 ENA(O)

B2 FLAG

B4 FLAG

B6 FLAG

Q-REG FLAG

ADDRESS FLAG

CLEAR CONTROL FLAGS
AND BREAKPOINT

STA(BIN)
 LDA(BK)
 STA(BCD)
 STA(UBCD)
 ENA(13101B)
 SSH(BREAK)
 STA(ERASE)
 OUT(ERASE)
 RAD(ERASE)
 ENA(15K1)
 EXF(77205B)
 EXF(77403B)
 EXF(77521B)
 EXF(77305B)
 EXF(77321B)
 EXF(77503B)
 EXF(77303B)
 EXF(77503B)
 EXF(77341B)
 EXF(77211B)
 EXF(77241B)
 EXF(77421B)
 SLJ(23K)

SSH(AFLG)
 SSH(PFLG)
 SSH(START)
 SSH(DUMP)
 SSH(BREAK)
 ALS(9)
 ENA(150B)
 ENA(2000B)
 JPI(14K1)
 SAU(14K1)
 ISK1(40)
 ISK1(40)
 ISK1(40)
 ISK1(40)
 ISK1(40)
 ISK1(40)
 ISK1(40)
 ISK1(40)
 ISK1(40)
 ISK1(40)
 ISK1(40)
 ZRD(0)

ERASE
 MEMORY
 DUMP
 DISPLAY

CLEAR
 LIGHTS

* KEYBOARD 1

EXF(777172B)
 EXF(77140B)
 SSK(BREAK)
 SSK(4K1)
 SLJ(4K1)
 SSK(B2FLG)
 SLJ(5K1)
 SSK(B3FLG)
 SLJ(6K1)
 SSK(8FLG)
 SLJ(7K1)
 SSK(85FLG)
 SLJ(8K1)
 SSK(B6FLG)
 SLJ(9K1)
 SSK(AFLG)
 SLJ(10K1)
 SSK(QFLG)
 SLJ(11K1)

SLJ(1K)
 INT(KYBD1)
 SLJ(3K1)
 SLJ(L+2)
 ZRD(0)
 SLJ(L+2)
 ZRD(0)
 SLJ(L+2)
 ZRD(0)
 SLJ(L+2)
 ZRD(0)
 SLJ(L+2)
 ZRD(0)
 SLJ(L+2)
 ZRD(0)
 SLJ(L+2)
 ZRD(0)
 SLJ(L+2)
 ZRD(0)

SENSE KYBD 1
 INPUT KYBD 1
 PROCESS KYBD 1
 HIT
 HIT

ENQ(7700B)
 QLS(6)
 MEQ(BK1)
 LDL(ZERO)
 TSK1(4)
 EXF(77010B)
 OUT(BK1)
 SLJ(1K)
 ENA(0)
 SLJ4(15K1)
 LDA(BIN)
 SLJ(1D+1)
 SLJ4(15K1)
 LDA(BIN)
 SLJ(1D+1)
 SLJ4(15K1)
 LDA(BIN)
 SLJ(1D+1)
 SLJ4(15K1)
 LDA(BIN)
 SLJ(1D+1)
 SLJ4(15K1)
 LDA(BIN)
 SLJ(1D+1)
 SLJ4(15K1)
 LDA(BIN)
 ENA(1D+6)
 ENQ(4)
 ENQ(KYB01)
 SSSH(UPLOW)
 SLJ4(25K)

3K1B
4K1

5K1

6K1

7K1

8K1

9K1

10K1

11K1

12K1
13K1

14K1

QLS(6)
 ENA(0)
 SLJ(L+2)
 RAD(BK1)
 SLJ(L-3)
 ENI(13)
 RSD(BRKDEX)
 ZRO(0)
 SLJ(3K1A)
 ZRO(0)
 SAU(DEX)
 ZRO(0)
 ZRO(0)
 SAL(DEX)
 ZRO(0)
 ZRO(0)
 SLJ(1D+1)
 ZRO(0)
 SLJ(1D+1)
 ZRO(0)
 SLJ(1D+1)
 ZRO(0)
 STA(TEMA)
 ZRO(0)
 ZRO(0)
 STA(TEHQ)
 ZRO(0)
 ZRO(0)
 ZRO(0)
 STA(ADDRESS)
 AJP3(L+2)
 ZRO(0)
 ZRO(0)
 SAU(L+1)
 ZRO(0)
 ZRO(0)
 STA(PREG)
 AJP3(L+2)
 ZRO(0)
 SAU(14K1)
 ZRO(0)
 SSSH(UPLOW)
 AJP1(L+2)
 ENI(0)
 ZRO(0)

REPLACE ZEROS
WITH THEIR
EQUIVALENT
BCD CODE

DISPLAY MODIFIED BRKPT

B1

B2

B3

B4

B5

B6

A-REG

Q-REG

FUNCTION CODE
ADDRESS

P-REG

IS UPPER OR LOWER ADDRESS DESIRED

SLJ(1A) ZR(0)

* FIFTEEN BIT REGISTERS

15K1 SLJ(N)
 AJP2(L+2)
 SLJ(1K)
 ADD(BRKDEX)
 SAL(L+4)
 ENQ(70B)
 AJP(15K1B)
 LDL(KYBD1)
 ALS(N)
 QLS(N)
 STA(BIN)
 SLJ(15K1)
 SLJ(15K1A)

* 48 BIT REGISTERS

16K1 SLJ(N)
 INAL(16)
 LDA(BRKDEX)
 ADD(BRKDEX)
 SAL(L+4)
 LDL(KYBD1)
 AJP(16K1B)
 LDL(KYBD1)
 ENQ(7770B)
 ADL(BIN)
 RAD(BRKDEX)
 ENAL(0)
 SLJ(23K)
 SLJ(1K)

* JUMP TABLE TABULATION

1T ENAL(1)
 AJP3(1A)
 SUB(TEMP)
 LDA(TEMP)
 LIL2(STORE)
 MEGS(BUF1)
 IJP2(4T)
 IJL2(STORE)
 LDA(ADDRESS)
 ENAL(1)
 LDAL(3)

2T
 3T

IF ABOVE OR BELOW
 MAIN PROGRAM
 DO NOT TABLE JUMP ADDRESSES

S A V E
 ADD JUMPED FROM
 ADD JUMPED TO
 NO. OF TIMES EXECUTED


```

LDA(L2)
LDA(L1)
SIL2(L1)
SLJ(5T)
INI2(1)
MEQ(ADDRESS)
RAD2(BUF3)
STA(L4)
STA(L3)
STA(L2)
STA(L1)
LNA(49)
MTH(STORE)
SLJ(1P)
ENA(2600B)
MEQ(PREG)
SLJ(1P)
LDA(TEMP)
STA(L3)
STA(L2)
INR(0)
LDA2(BUF2)
SLJ(7T)
LDA(L2)
LDA(L1)
SIL2(L1)
SLJ(6T)
ENI(0)
SLJ(6T)
ENI(0)
ENQ(37600B)
SLJ(1A)
ENI(0)
SLJ(2T)

```

```

IF BUF FULL
GO PRINT
IF NOT IN RESIDENT GO TO 1A
IF IN RESIDENT GO PRINT

```

* PRINT JUMP TABLE

```

1P LDA(PRTJUM)
SIL2(NAN)
PRINT 10
FORMAT(//76X,22HFROM TO
10 FORMAT(1,(BUF1(I),BUF2(I),BUF3(I),I=1,NAN)
11 FORMAT(5X,05,4X,05)
2P MEQ(2600B)
ENQ(37600B)
SLJ(3P)
ENI(0)
SLJ(L+4)
OUT1(END)
ENI(1D)
SAU(0)
SAU(L+3)
LIL2(DEMA)
LDA(TEMP)
ZRO(0)
ENI(0)
STA2(BUF2)
ENI(0)
SLJ(L-2)
SLJ(LA)
ENI(0)

```

• CHECK PRINT FLAG

RELINQUISH CONTROL

TO MONITOR
CLEAR BUFS

* ERASE DD65 AND DISPLAY INTRODUCTION

```

1E SLJ(N)
ENI(0)

```


thesL636

On-line program analysis and diagnosis.



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