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Artificial Intelligence through Prolog by Neil C. Rowe

Rowe, Neil C.

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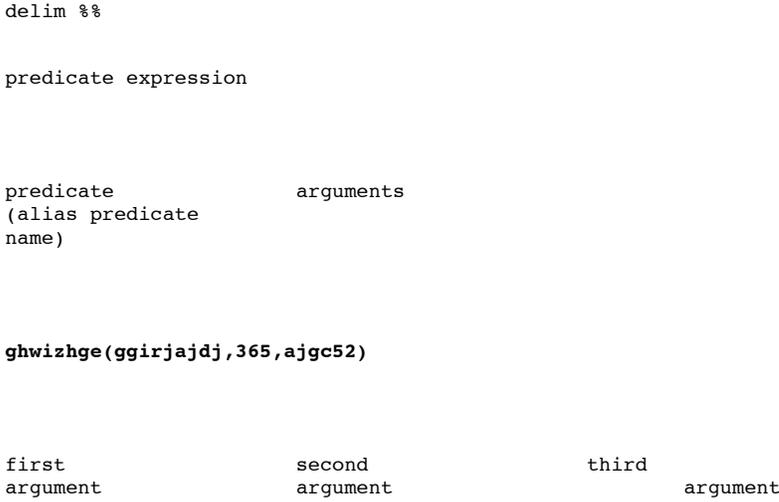
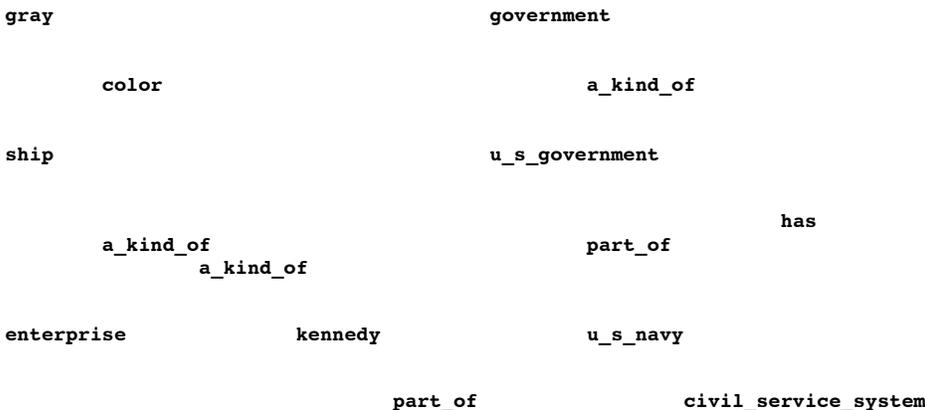


Figure 2-1: terminology about predicate expressions



Figure 2-2: summary of the kinds of predicates in Chapter 2



part_of

location

15n35e

Figure 2-3: an example semantic network

Symbol	Description	Meaning
?-	question mark,	query prompt
	minus sign	(typed by Prolog)
.	period	(1) query terminator (typed by user)
		(2) database fact and rule terminator
,	comma	(1) "and" of predicate expressions in a query
		(2) separator of arguments
		(3) separator of list items (see chapter 5)
;	semicolon	(1) "or" of predicate expressions in a query
		(2) forces new alternative after a query answer
()	parentheses	(1) for grouping arguments together
		(2) for grouping query expressions
_	underscore	no special meaning, but we often use
	as a character	to make long names legible
:-	colon,	rule definition symbol (see
	minus sign	chapter 4); read as "if"
[]	square brackets	a list (see chapter 5)
<>	angular brackets	not a Prolog symbol, but
		we use it to enclose descriptions of
		things (like Backus-Naur Form)

Figure 3-2: special symbols used with Prolog in this book

Query:

```
?-boss(X,Y), boss(Y,Z).
```

Database:

```
boss(dick,harry).
boss(tom,dick).
boss(ann,mary).
boss(mary,harry).
```

Processing:

```
Step number/First boss predicate expression:/Second boss predicate expression:
in text/Bindings made/Binding of Z
/of X and Y

1/X=dick,Y=harry/
2/|[no Z possible]
3/X=tom,Y=dick/
4/|Z=harry
5/|
```

Figure 3-3: processing state at generation of first answer to superboss query: X=tom,Y=dick,Z=harry

Query:

```
?-boss(X,Y), boss(Y,Z).
```

Database:

```
boss(dick,harry).
boss(tom,dick).
boss(ann,mary).
boss(mary,harry).
```

Processing:

```

Step number/First boss predicate expression:/Second boss predicate expression:
in text/Bindings made/Binding of Z
/of X and Y
1/X=dick,Y=harry/
2|[no Z possible]
3/X=tom,Y=dick/
4|[Z=harry
5|/
6|/
7|[no further Z possible]
8/X=ann,Y=mary/
9|[Z=harry/
10|/

```

Figure 3-4: processing state at generation of second answer to superboss query: X=ann,Y=mary,Z=harry

Query:

```
?-boss(X,Y), not(boss(X,dick)), boss(Y,Z).
```

Database:

```

boss(dick,harry).
boss(tom,dick).
boss(ann,mary).
boss(mary,harry).

```

Processing:

```

Step number/First predicate/Second predicate/Third predicate
in text/expression:/expression:/expression:
/Bindings made/What happens?/Binding of Z
/of X and Y
1/X=dick,Y=harry
2|succeeds
3|[no possible Z]
4|fails
4/X=tom,Y=dick
5|fails/
6/X=ann,Y=mary
7|succeeds
8|[Z=harry

```

Figure 3-5: processing order for the extended superboss example: answer X=ann,Y=mary,Z=harry

Database:

```

Memory location/Contents
10000|a_kind_of(enterprise,ship).
10026|a_kind_of(kennedy,ship).
10052|part_of(enterprise,u_s_navy).
10070|part_of(kennedy,u_s_navy).
10088|part_of(u_s_navy,u_s_government).

```

Index:

Predicate name/Locations of facts using it

```

a_kind_of|10000,10026
part_of|10052,10070,10088

```

Figure 3-6: an example of predicate indexing

left side	special symbol	right side
(a predicate	meaning "if"	(a query)
expression)		

```
a(X,middle,Y) :- b(X), c(X,Y,extra), d(X,Z), e(Z,Y), f(Z).
```

parameter variables local variable rule terminator

Read a rule as meaning: the left side is true if querying the right side succeeds. Figure 4-1: terminology about rules

Query:

```
?-boss(X,Y), not(boss(X,tom)).
```

Database:

```
department(tom,sales).
department(harry,production).
manager(dick,sales).
manager(mary,production).
boss(B,E) :- department(E,D), manager(B,D).
```

Processing:

```
Step/In rule/In rule/In query:/In rule/In rule/In query:
in text/call 1:/call 1:/Bindings/call 2:/call 2:/argument
/Bindings/Binding/to X and Y/Bindings/Binding/to not
/to E and D/to B/to B and E/to D/
```

```
1|E=tom,
|D=sales
2||B=dick/
2|||X=dick,
|||Y=tom
3|||/
4|||B=dick,
|||E=tom
4|||D=sales
4|||/succeeds
5||[no more
||choices]
6|E=harry,
|D=production
7||B=mary
7||X=mary,
||Y=harry
8|||B=mary,
|||E=tom
9|||D=sales
10|||/fails
```

Figure 4-2: processing order of the other-than-tom's-boss query: answer X=mary,Y=harry

```
X      r      Z      r      Y
      r
```

Figure 4-3: the general form of transitivity

```
d
c
b
a
```

```
on(b,a).
on(c,b).
on(d,c).
```

Figure 4-4: a transitivity example

y	p	Value
r		
x	p	

Figure 4-5: the general form of inheritance

truck_4359	location	nps_north_lot
	owner	nps
part_of	contains	status
		status defective
battery_of_truck_4359	location	
	owner	
	status	defective

Figure 4-6: an inheritance example

vehicle	purpose	transportation
	a_kind_of	
ship		
	a_kind_of	
carrier		a_kind_of
	a_kind_of	
vinson	purpose	

Figure 4-7: an example of both transitivity and inheritance

	action		
		pedsignals	
safe_stop_possible	pedhalt	pedgo	greenfacing

light

clockwise_cross

Figure 4-8: the predicate hierarchy for the traffic lights program

```

D
C   G
B   F   I
A   E   H
-----

```

Figure 4-9: some blocks on a table

Predicate/Arguments/What it means

```

first(L,I) | a list, | I is the first
            | an item | item of the list L
last(L,I)  | a list, | I is the last
            | an item | item of the list L
member(I,L) | an item, | I is an item somewhere
            | a list  | inside the list L
length(L,N) | a list, | N is the length
            | a number | of the list L
max(L,N)    | a list, | N is the largest number
            | a number | in the list of numbers L
delete(I,L1,L2) | an item, | L2 is L1 with all
                | a list,  | occurrences of the I
                | a list  | removed from it
append(L1,L2,L3) | a list, | L3 is the glueing
                 | a list, | together of the
                 | a list  | lists L1 and L2
subset(L1,L2)   | a list, | all items in L1
                 | a list, | are in L2
sort(L1,L2)     | a list, | L2 is the sorting
                 | a list  | of list of numbers L1

```

Figure 5-1: the classic list-processing predicates defined in this chapter (to be used in the rest of this book)

Given the definition of `append`:

```

append([],L,L).
append([X|L],L2,[X|L3]) :- append(L,L2,L3).

```

Given the query:

```
?- append([gas,oil],[tires,battery,radiator],Things_to_check_before_trip).
```

Nested environments created:

```

Query: ?- append([gas,oil],[tires,battery,radiator],Things_to_check_before_trip).
In definition line 2: [X|L]=[gas,oil], L2=[tires,battery,radiator], [X|L3]=???
So X=gas, L=[oil], [X|L3]=[gas|???]

```

```

Query: ?- append([oil],[tires,battery,radiator],L3).
In definition line 2: [X|L]=[oil], L2=[tires,battery,radiator], [X|L3]=???
So X=oil, L=[], [X|L3]=[oil|L3]

```

```

Query: ?- append([], [tires,battery,radiator],L3).
In definition line 1: L=[tires,battery,radiator]
Query succeeds

```

Note: "???" means unknown Figure 5-2: an example of using `append` with first two arguments bound (inputs) and last argument unbound (an output); the processing state is just after the innermost recursion succeeds

Given the definition of `append`:

```

append([],L,L).
append([X|L1],L2,[X|L3]) :- append(L1,L2,L3).

first/second/third/What happens?/Example
arg./arg./arg./

bound|bound|bound|Answers yes if third|?- append([a,b],[c],[a,b,c]).
||argument is the first two|yes.
||arguments glued together
||in order), else no
bound|bound|unbound|Glues the second argument|?- append([a,b],[c,d],L).
||on the end of the first|L = [a,b,c,d]
||argument, binds result
||to the third argument
bound|unbound|bound|Checks if the first|?- append([a,b],L,[a,b,c]).
||argument is on the,|L = [c]
||front of the third,|
||binding the second to
||the rest of the third
unbound|bound|bound|Checks if the second|?- append(L,[c,d],[a,b,c,d]).
||argument is on the|L = [a,b]
||back of the third,
||binding the first to
||the rest of the third
unbound|unbound|bound|Generates all divisions|?- append(L1,L2,[a,b,c]).
||of the third argument|L1 = [], L2 = [a,b,c];
||into two pieces|L1 = [a], L2 = [b,c];
||preserving term order|L1 = [a,b], L2 = [c];
||L1 = [a,b,c], L2 = []
unbound|bound|unbound|Generates in abstract|?- append(L1,[a,b],L2).
||form all results of|L1 = [], L2 = [a,b];
||glueing something on|L1 = [_1], L2 = [_1,a,b];
||the front of the|L1 = [_1,_2], L2 = [_1,_2,a,b]
||second argument|(The _1 and _2
||are Prolog-invented
||variable names.)
bound|unbound|unbound|Generates in abstract|?- append([a,b],L1,L2).
||form all results of|L1 = _1, L2 = [a,b|_1]
||glueing something on|(The _1 is a
||the back of the|Prolog-invented
||first argument|variable name.)

```

Note: any bound arguments to append must be lists for the above program to work.

Figure 5-3: the different possible uses of the append predicate

*increasing
flexibility*

meta-rules

backwards
chaining

forwards
chaining

hybrid
chaining

virtual facts

partitioning

parallelism

and-or-not lattices

decision lattices

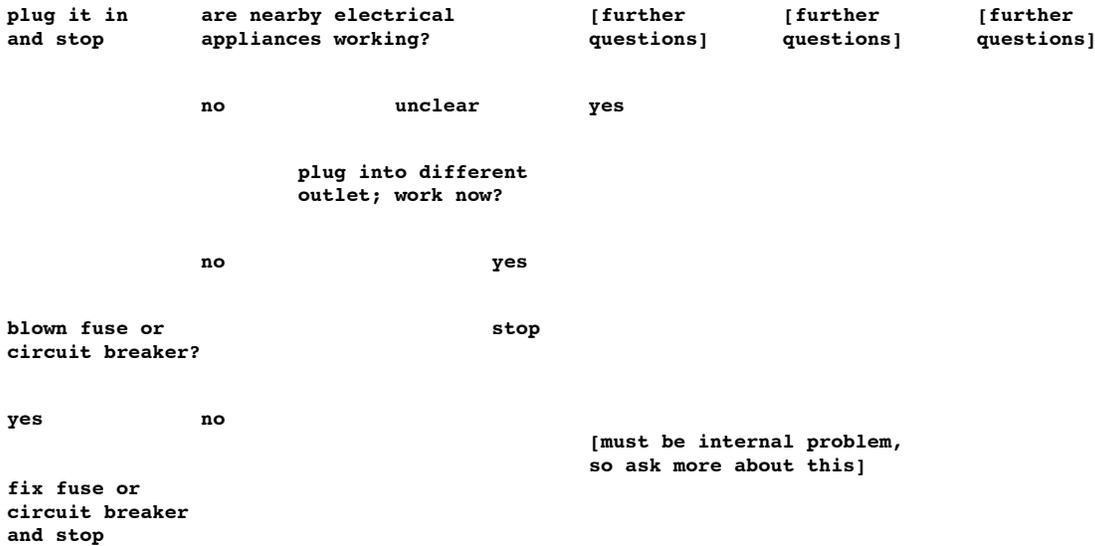


Figure 6-5: a decision lattice

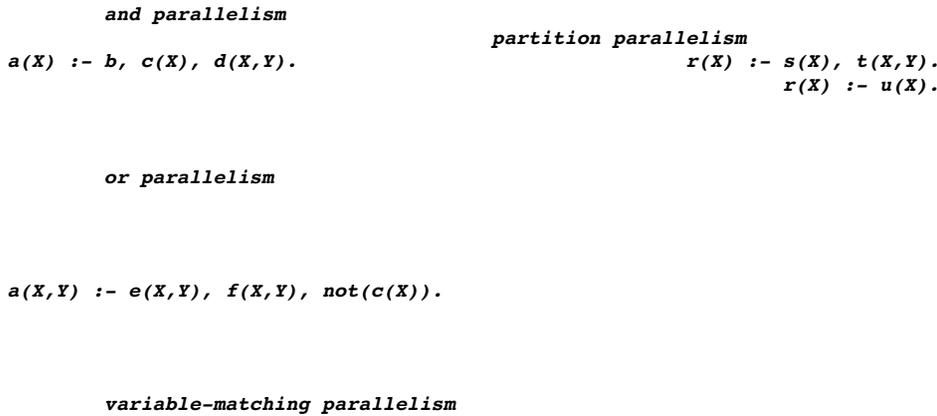
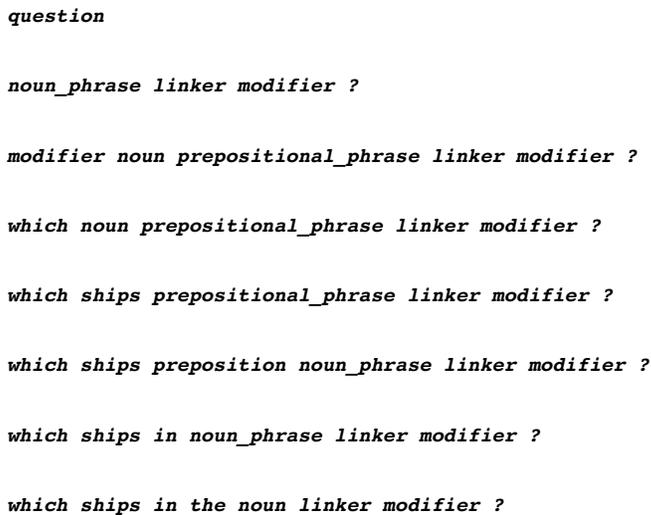


Figure 6-6: types of parallelism



which ships in the Mediterranean linker modifier ?

which ships in the Mediterranean are modifier ?

which ships in the Mediterranean are American ?

Figure 6-8: a parsing example (downwards arrows represent top-down parsing, upwards arrows represent bottom-up)

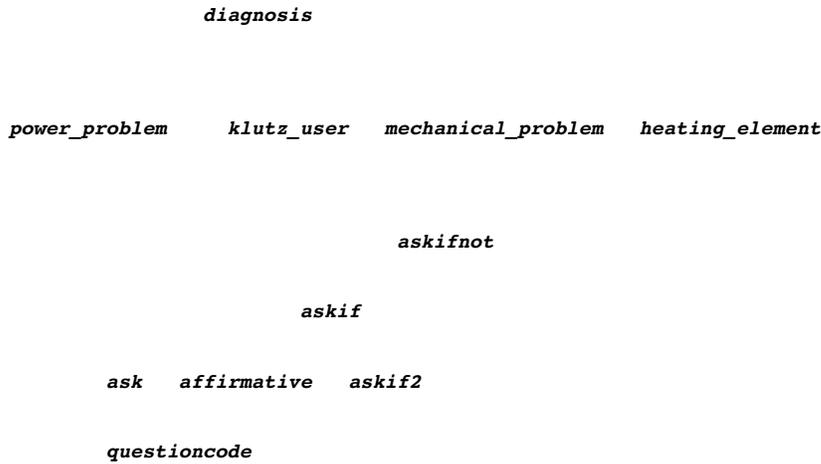


Figure 7-2: the predicate hierarchy for the appliance diagnosis expert system, including the problem-independent predicates

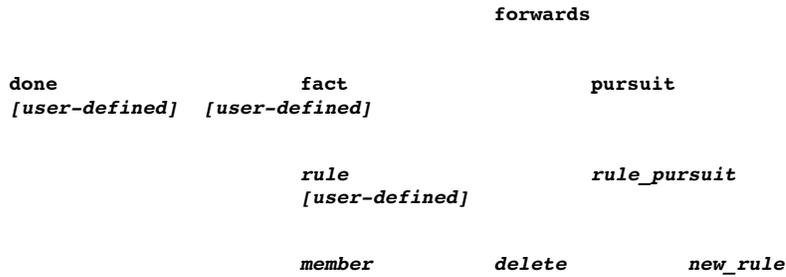
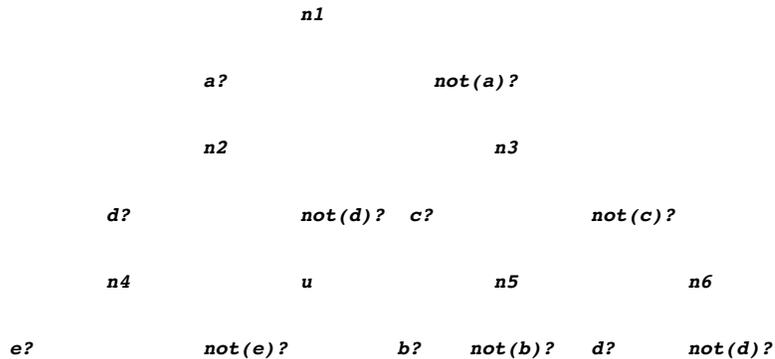


Figure 7-3: the predicate hierarchy for the forwards-chaining program



u r t v t s

Figure 7-4: a derived decision tree

Independence-assumption "or":
 $p(A \text{ or } B) = p(A) + p(B) - p(A)p(B)$
 Independence-assumption "and":
 $p(A \text{ and } B) = p(A)p(B)$

 Conservative "or":
 $p(A \text{ or } B) = p(A) \text{ if } p(A) > p(B), \text{ else } p(B)$
 Liberal "and":
 $p(A \text{ and } B) = p(B) \text{ if } p(A) > p(B), \text{ else } p(A)$

 Liberal "or":
 $p(A \text{ or } B) = p(A) + p(B) \text{ if this } < 1, \text{ else } 1$
 Conservative "and":
 $p(A \text{ and } B) = p(A) + p(B) - 1 \text{ if this } > 0, \text{ else } 0$

p(A) p(B)

Figure 8-2: Venn diagrams illustrating the three standard probability-combination methods, applied to two probabilities

/and-combination/or-combination

Independence assumption | % p sub 1 p sub 2 ... p sub n | % $1 - [(1 - p_{sub 1})(1 - p_{sub 2}) \dots (1 - p_{sub n})]$
 (reasonable-guess)
 Conservative assumption | % max (0 , (p sub 1 + p sub 2 + ... + p sub n) - n + 1) | % max (p sub 1 , p sub 2 , ... p sub n) | %
 (lower bound)
 Liberal assumption | % min (p sub 1 , p sub 2 , ... p sub n) | % min (1 , p sub 1 + p sub 2 + ... + p sub n) | %
 (upper bound)

Note:
 conservative "and" % <= % independence-assumption "and" % <= % liberal "and"
 % <= % conservative "or" % <= % independence-assumption "or" % <= % liberal "or"

Figure 8-3: formulas for combination of probabilities %p sub i%, i=1 to n

|definitely|probably|probably not|definitely not|
 definitely|definitely|definitely|definitely|definitely
 probably|definitely|probably|probably|probably
 probably not|definitely|probably|probably not|probably not|
 definitely not|definitely|probably|probably not|definitely not|

Figure 8-5: one evidence combination method

|definitely|probably|probably not|definitely not|
 definitely|definitely|definitely|definitely|definitely
 probably|definitely|definitely|definitely|probably
 probably not|definitely|definitely|probably|probably not|
 definitely not|definitely|probably|probably not|definitely not|

Figure 8-6: another evidence combination method

a b c e f

d
g h

Figure 9-2: the equivalent search graph for Figure 9-1

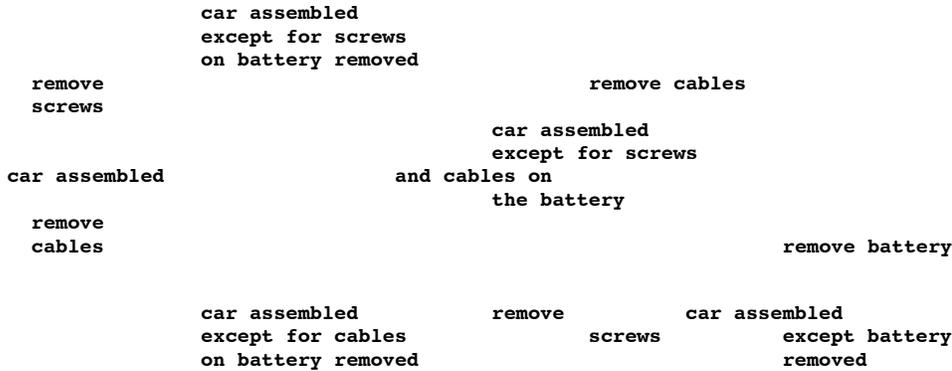


Figure 9-3: part of a search graph for auto repair

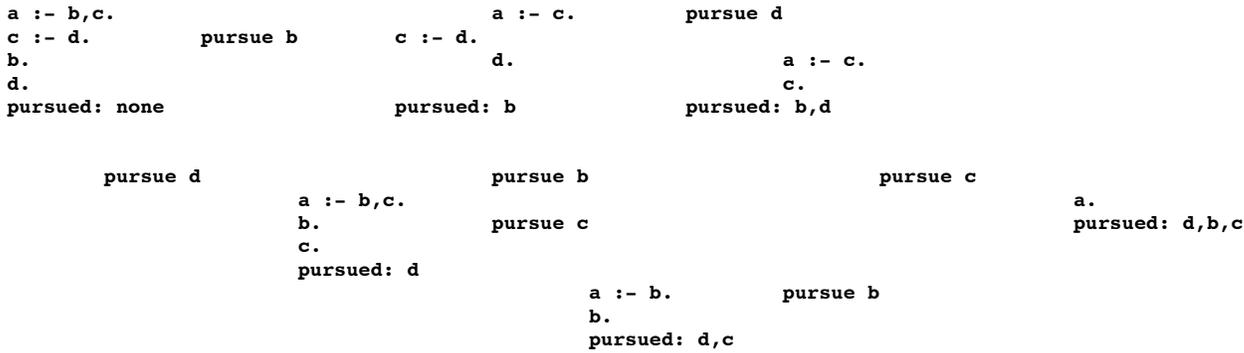


Figure 9-4: a complete search graph for generalized (any-fact-order) forwards chaining from a particular set of facts and rules (rules no longer useful are deleted from states as search proceeds)

Name of/Uses/Uses/Uses/Next state whose
search strategy/agenda?/evaluation/cost/successors are found
//function?/function?

Depth-first|no|no|no|A successor of the last
search|||state, else the predecessor
Breadth-first|yes|no|no|The state on the
search|||agenda the longest
Hill-climbing|no|yes|no|The lowest-evaluation
(optimization)|||successor of the last state
Best-first|yes|yes|no|The state on the agenda
search|||of lowest evaluation value
Branch-and-bound|yes|no|yes|The state on the agenda
|||of lowest total cost
A* search|yes|yes|yes|The state on the agenda
|||of lowest sum of evaluation
|||value and total cost

Figure 9-5: the classic search strategies (heuristics may be used with any of these)

Given the rule set:

```
t :- a, b.
t :- c.
u :- not(a), c.
u :- a, d.
v :- b, e.
```

Then this is a decision lattice implementing it:

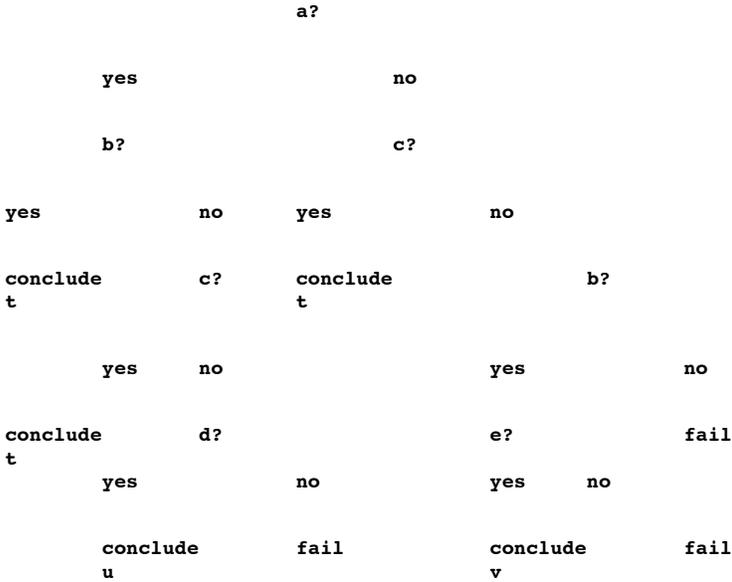


Figure 9-10: the decision lattice for backwards chaining on an example rule set

Assume the rules:

```
t :- a, b.
t :- c.
u :- not(a), c.
u :- a, d.
v :- b, e.
```

And assume the only possible facts are a, b, c, d, and e, provided for a database in order of priority a, c, b, d, and e. Assume each has an independent probability of occurrence P. Situations to consider:

a?	c?	b?	d?	e?	probability	number of matches	conclusion
false	true	-	-	-	$P (1 - P)$	1	t
true	true	-	-	-	P^2	3	t
true	false	true	-	-	$P^2 (1 - P)$	3	t
true	false	false	true	-	$P^2 (1 - P)$	2	u
true	false	false	false	true	$P^2 (1 - P)$	3	-
true	false	false	false	false	$P^2 (1 - P)$	4	-
false	false	true	true	true	$P^3 (1 - P)$	2	v
false	false	true	true	false	$P^3 (1 - P)$	3	-
false	false	true	false	true	$P^3 (1 - P)$	3	v
false	false	true	false	false	$P^3 (1 - P)$	4	-
false	false	false	true	true	$P^3 (1 - P)$	3	-
false	false	false	true	false	$P^3 (1 - P)$	4	-
false	false	false	false	true	$P^3 (1 - P)$	4	-
false	false	false	false	false	$P^3 (1 - P)$	5	-

Figure 9-11: analysis of cost of forwards chaining on the rule set

Given this array of edgeness measures for a picture (array g(i,j)):

0	1	0	8	2
1	2	7	3	1
3	1	5	2	0
2	2	3	3	1

```

2|5|2|1|3
1|7|1|2|1
3|6|2|2|1

```

Inference: there's a continuous edge running north-northeast to south-southwest, though it's hard to see towards the center of the picture.
That's the interpretation (array e(i,j)):

```

false|false|false|true|false
false|false|true|false|false
false|false|true|false|false
false|false|true|false|false
false|true|false|false|false
false|true|false|false|false
false|true|false|false|false

```

Figure 9-13: example arrays for finding visual edges

```

          b          c
a
          d          e

```

Figure 10-3: a simpler route-planning problem

Step number/depthsearch2/depthsearch2/depthsearch2/depthsearch2
in text|first call|second call|third call|fourth call

```

1|called with
|state a
2|first rule
|fails
3|second|called with
|rule tried|state b
4| |first rule|called with
||fails,|state c
||second tried
5| |both rules
||fail on c
6|backtrack to
|not, then to
|successor;
|choose d
7| |called with
||state d;
||first rule
||fails; 2nd
||rule picks
||b, which
||fails not
8| |choose e|called with
||state e
9| |first rule
||succeeds
||with path
||list [e,d,b,a]
||succeeds
|succeeds
|succeeds

```

Figure 10-4: summary of the depth-first search example

```

          breadthsearch
find_successors          cleandatabase          measurework

```



Figure 10-5: the rule-predicate hierarchy for the breadth-first search (breadthsearch) program

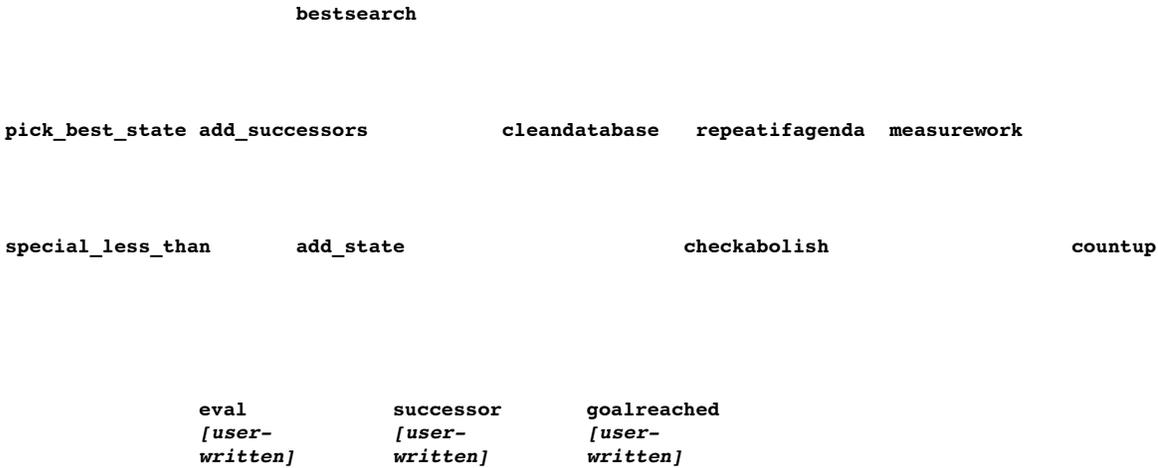


Figure 10-6: the rule-predicate hierarchy for the best-first search (bestsearch) program



Figure 10-7: the rule-predicate hierarchy for the A* search (astarsearch) program

```

Desired/Should/Should/Should/Should/Should
situation|you|you|you|you|you
|visit|visit|visit|visit|visit
|cafeteria?|vending|your|secretary?|colleague?
||machine?|office?
  
```

```

You want|yes|yes|no|no|no
to be not|
hungry,
and you're
hungry now
You want|no|no|yes|yes|no
to have
change,
and you
don't now
You want|no|no|yes|yes|yes
to know|
where
something
is

```

Figure 11-1: tabular representation of the operator recommendations (difference table) for the office-worker problem

```

Desired/Should/Should/Should/Should/Should
situation|you|you|you|you|you
|visit|visit|visit|visit|visit
|vending|cafeteria?|colleague?|secretary?|your|
|machine?|/|/|/|office?

```

```

You want|yes|yes|no|no|no
to be not|
hungry,
and you're
hungry now
You want|no|no|no|yes|yes
to have
change,
and you
don't now
You want|no|no|yes|yes|yes
to know|/|/|(but try)|(but try
where|/|/|third)|second)
something
is

```

Figure 11-2: a different tabular representation of the operator recommendations (difference table) for the office-worker problem

```

Desired/Should/Should/Should/Should/Should/Should/Should/Should
situation|you|you|you|you|you|you|you|you
|replace|replace|disassemble|disassemble|assemble|assemble|turn over|smash
|batteries|light|case|top|case|top|case|case
|?/?/?/?/?/?/?/?/?

```

```

You want|yes|no|no|no|no|no|no|no
batteries
OK, when
they aren't
You want|no|yes|no|no|no|no|no|no
light OK,
when it
isn't
You want|no|no|yes|no|no|no|no|yes
case open,
when it
isn't
You want|no|no|no|yes|no|no|no|yes
top open,
when it
isn't
You want|no|no|no|no|yes|no|no|no
case closed,
when it
isn't
You want|no|no|no|no|no|yes|no|no
top closed,
when it
isn't
You want|no|no|no|no|no|no|yes|yes
batteries

```

outside,
when they
aren't

Figure 11-4: tabular representation of the operator recommendations (difference table) for the flashlight problem

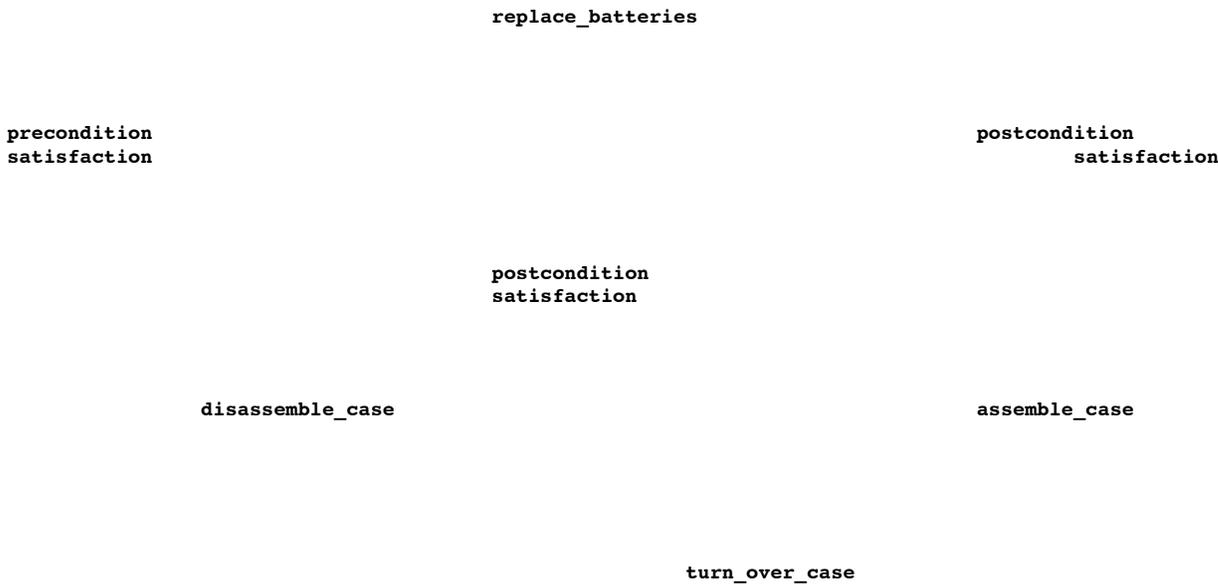


Figure 11-6: the form of the solution of the dead-batteries flashlight problem by means-ends analysis (Figure 11-7 gives a more detailed picture)

.DS L Level 1 State=[closed(case),closed(top),inside(batteries),defective(batteries), ok(light),unbroken(case)] Goal=[ok(batteries),closed(case),closed(top)] Oplist=[disassemble_case,turn_over_case,replace_batteries,assemble_case] Goalstate=[closed(case),inside(batteries),ok(batteries),closed(top), ok(light),unbroken(case)] Operator=replace_batteries Level 2 (Precondition recursion) (Postcondition recursion) State=[closed(case),closed(top), State=[inside(batteries),ok(batteries), inside(batteries),defective(batteries), open(case),closed(top), ok(light),unbroken(case)] ok(light),unbroken(case)] Goal=[open(case),outside(batteries),] Goal=[ok(batteries),closed(case), unbroken(case)] closed(top)] Oplist=[disassemble_case,turn_over_case] Oplist=[assemble_case] Goalstate=[outside(batteries),open(case), Goalstate=[closed(case),inside(batteries), closed(top),defective(batteries), ok(batteries),closed(top), ok(light),unbroken(case)] ok(light),unbroken(case)] Operator=disassemble_case Operator=assemble_case Level 3 (Precondition (Postcondition (Precondition (Postcondition recursion) recursion) recursion) recursion) [nothing State=[open(case),closed(top), [nothing [nothing needed] inside(batteries),defective(batteries) needed] needed] ok(light),unbroken(case)] Goal=[outside(batteries)] Oplist=[turn_over_case] Goalstate=[outside(batteries),open(case), closed(top),defective(batteries), ok(light),unbroken(case)] Operator=turn_over_case Level 4 (Precondition recursion) (Postcondition recursion) [nothing needed] [nothing needed]

Figure 11-7: the recursive calls in the dead-batteries flashlight problem, with bindings eventually found for Oplist, Goalstate, and Operator

```
frame_name: business_form
a_kind_of: physical_object
type:
author:
```

```
frame_name: memo
a_kind_of: business_form
type: informal
author:
addressees:
subject:
date:
text:
```

```

frame_name: memos of Ann
a_kind_of: memo
type: informal
author: Ann (D)
addressees:
subject:
date:
text:

frame_name: memos to Ann
a_kind_of: memo
type: informal
author:
addressees: [Ann] (D)
subject:
date:
text:

frame_name: budget memos
a_kind_of: memo
type: informal
author:
addressees:
subject: budget (D)
date:
text:

frame_name: budget memos to Ann from Tom
a_kind_of: memos to Ann
a_kind_of: budget memos
type: informal
author: Tom
addressees: [Ann] (D)
subject: budget (D)
date:
text:

frame_name: budget memo to Ann from Tom on 6/12
a_kind_of: budget memos to Ann from Tom
type: informal
author: Tom (D)
addressees: [Ann] (D)
subject: budget (D)
date: 6/12 (D)
text: "Error in personnel estimate:
      10, not 10,000."

```

Notes: each box is a frame, and arrows represent a_kind_of facts. Slot names are given, then a colon, and then the slot value; a "(D)" means the value follows by the definition of the frame.

Figure 12-1: some frames about memos

```

ship      part_of      hull

a_kind_of

enterprise      part_of      enterprise_hull

a_kind_of

```

Figure 12-2: an example of part-kind inheritance

```

physical_object

a_kind_of

vehicle

a_kind_of

car      extension      cars_now_on_the_road

part_of

a_kind_of

```

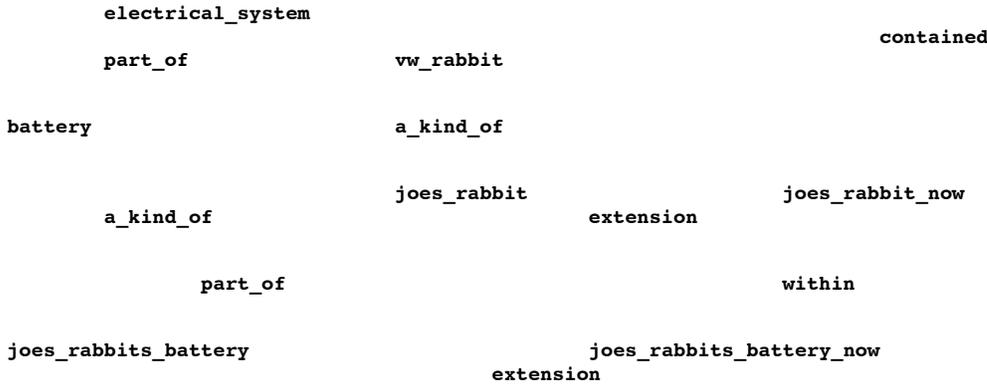


Figure 12-3: the semantic network for the frames in the cars example

```

frame_name: sending8347
a_kind_of: sending
actor: we
object: memo72185
time: yesterday
destination: headquarters
method: express mail
  
```

```

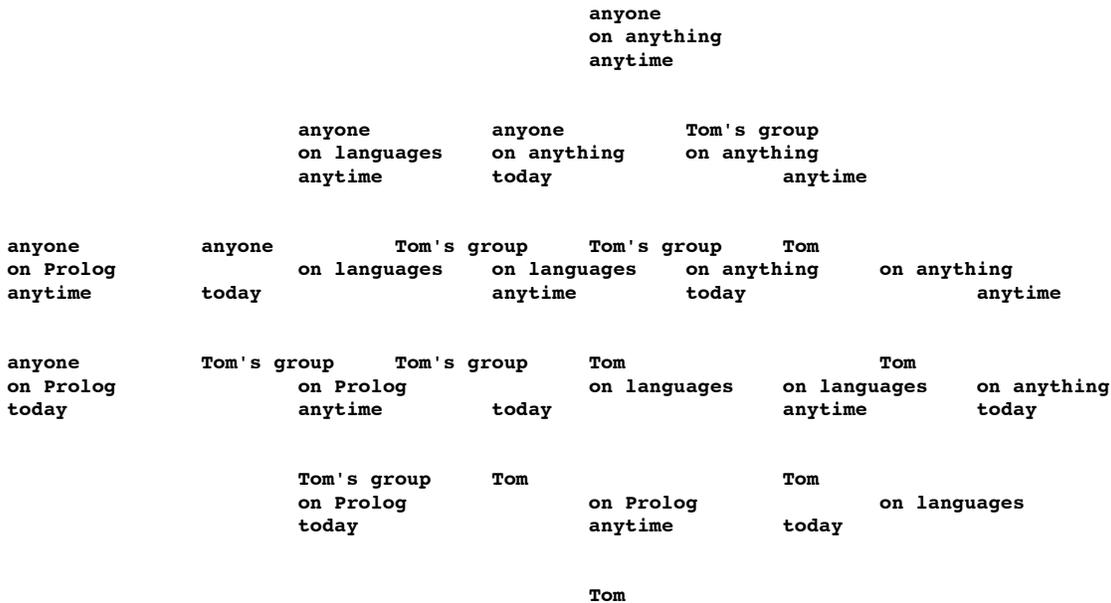
frame_name: memo72185
a_kind_of: memo
author: Tom
addressees: [Ann]
origin: drawingup20991
  
```

```

frame_name: drawingup20991
a_kind_of: drawingup
actor: Tom
object: memo72185
beneficiary: Ann
time: 6/12
  
```

.sp 7

Figure 12-4: frame representation of the meaning (semantics) of the sentence "Yesterday we sent headquarters by express mail the budget memo that Tom drew up for Ann on 6/12"



on Prolog
today

Note: "languages" means "programming languages", "Prolog" means "Prolog interpreter"; all arrows represent a_kind_of

Figure 12-5: a three-dimensional inheritance hierarchy of user model frames

```
Time/Monday/Tuesday/Wednesday/Thursday/Friday
9|occupied|occupied|occupied||occupied
10|occupied|occupied|occupied|occupied|occupied
11|class|occupied|occupied|class|occupied
12|occupied|occupied|occupied|occupied|occupied
1|class|class|occupied|class|occupied
2|occupied||occupied||occupied
3|occupied|occupied|occupied|occupied|occupied
4|occupied|occupied|occupied||occupied
```

Figure 13-1: an example scheduling problem, with an example solution symbolized by "class"

```
value/Satisfies/Satisfies/Satisfies/Satisfies/Satisfies
/label?|large?|regular?|irregular?

grass|yes|yes|no|yes
water|yes|yes|no|yes
pavement|yes|yes|yes|no
house|yes|yes|yes|no
vehicle|yes|no|yes|no
```

Figure 13-3: summary of the single-argument predicates in the photo interpretation example

```
A1/A2/Satisfies/Satisfies
value/value/borders(A1,A2)?|inside(A1,A2)?

grass|grass|no|no
grass|water|yes|yes
grass|pavement|yes|yes
grass|house|yes|no
grass|vehicle|no|no
water|grass|yes|yes
water|water|no|no
water|pavement|yes|yes
water|house|no|no
water|vehicle|no|no
pavement|grass|yes|no
pavement|water|yes|no
pavement|pavement|no|no
pavement|house|yes|no
pavement|vehicle|yes|no
house|grass|yes|yes
house|water|no|no
house|pavement|yes|yes
house|house|no|no
house|vehicle|no|no
vehicle|grass|no|no
vehicle|water|no|no
vehicle|pavement|yes|yes
vehicle|house|no|no
vehicle|vehicle|no|no
```

Figure 13-4: summary of the two-argument predicates in the photo interpretation example

Given the query:

```
?- a(X), b(X,Y), c(Z), d(Y,Z), e(X).
```

These are its dependencies:

*Predicate/Variables/Dependencies to/Dependencies that
expression/bound/previous predicate/bind a common variable
//expressions*

```
a(X) | X | none | none
b(X,Y) | Y | a(X) | a(X)
c(Z) | Z | none | none
d(Y,Z) | none | b(X,Y), c(Z) | b(X,Y), c(Z)
e(X) | none | a(X), b(X,Y) | a(X)
```

Figure 13-5: summary of dependencies in an example query

Given the query:

```
?- a(X), b(X,Y), c(Z), d(X,Z), e(X).
```

Given the database:

```
a(1).
a(2).
a(3).
b(A,B) :- B is 3*A.
c(4).
c(1).
d(A,B) :- A>B.
e(3).
```

This is what happens in dependency-based backtracking:

```
Step|a(X)|b(X,Y)|c(Z)|d(X,Z)|e(X)

start|active|active|active|active|active
1|X bound to 1|active|active|active|active
2|inactive|Y bound to 3|active|active|active
3|inactive|inactive|Z bound to 4|active|active
4|inactive|inactive|inactive|fails|active
5|active|inactive|Z bound to 1|active|active
6|active|inactive|inactive|fails|active
7|active|inactive|fails|active|active
8|X bound to 2|active|active|active|active
9|inactive|Y bound to 6|active|active|active
10|inactive|inactive|Z bound to 4|active|active
11|inactive|inactive|inactive|fails|active
12|active|inactive|Z bound to 1|active|active
13|active|inactive|inactive|succeeds|active
14|active|inactive|inactive|inactive|fails
15|X bound to 3|active|inactive|active|active
16|inactive|Y bound to 9|inactive|active|active
17|inactive|inactive|inactive|succeeds|active
18|inactive|inactive|inactive|inactive|succeeds
```

Figure 13-6: a dependency-based backtracking example

n/p/Probability

```
5|.25|.76
10|.25|.94
20|.25|1.0E+00
50|.25|1.0E+00
5|.1|.4
10|.1|.65
20|.1|.88
50|.1|.99
5|.01|.05
10|.01|.10
20|.01|.18
50|.01|.39
100|.01|.63
```

Figure 13-7: the probability that some item in a possibility list will be impossible, for an n-item list with probability p that a possibility will be impossible

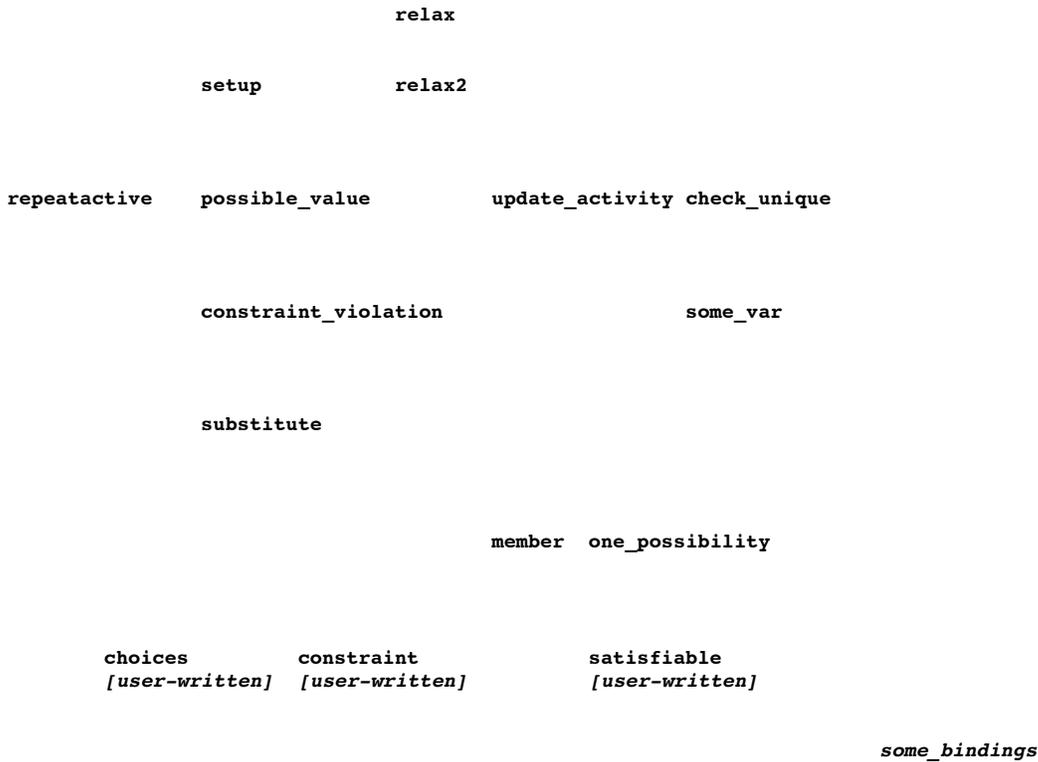


Figure 13-8: the predicate hierarchy for the pure-relaxation program

```

a/b/a; not(b) | Is the rule | Justification
||| a :- b. consistent
||| (uncontradicted)?
true|true|true|yes|b proves a
true|false|true|yes|a could be proved
||| by another rule, and
||| that wouldn't violate
||| the truth of this one
false|true|false|no|b proves a by the
||| rule, a contradiction
false|false|true|yes|the rule doesn't
||| apply, and there may
||| be no other rule to
||| prove a

```

Figure 14-1: demonstration that an equivalent logical (declarative) meaning of the Prolog rule a :- b. is a; not(b).

```

a; not(b); not(c).      c; not(d).      b.      d.

                                not(a).
                                [from step 1]

```

not(b); not(c).
[from step 2]

not(c).
[from step 3]

not(d).
[from step 4]

"null clause"
[from step 5]

Figure 14-2: a resolution example

In rule form/In clause form

BACKWARDS CHAINING:
QUERY: ?- a. | **ASSUMPTION:** not(a).
a :- b, c, d. | a; not(b); not(c); not(d).
RESULTS IN: | **RESOLVES TO:**
?-b, c, d. | not(b); not(c); not(d).
FORWARDS CHAINING:
a :- b, c. | a; not(b); not(c).
c. | c.
RESULTS IN: | **RESOLVES TO:**
a :- b. | a; not(b).
RULE COLLAPSING:
a :- b, c, d. | a; not(b); not(c); not(d).
c :- e, f. | c; not(e); not(f).
RESULTS IN: | **RESOLVES TO:**
a :- b, e, f, d. | a; not(b); not(e); not(f); not(d).

Figure 14-3: three special cases of resolution

1. methods requiring a gold standard

- (a) traces
- (b) confusion matrices
- (c) comparison of a near miss and example

- (i) no numbers involved
- (ii) composite results
- (iii) numbers involved

- (d) comparison of two examples

- (i) no numbers involved
- (ii) composite results
- (iii) numbers involved

2. methods not requiring a gold standard

- (a) schema for rules and facts
- (b) user-initiated debugging questions
- (c) cooperativeness evaluation
- (d) a priori suitability for artificial intelligence techniques

Figure 15-1: tools for testing and debugging

guessed/battery/electrical/fuel/engine
cause/dead/short/blockage/problem

```

actual
cause
battery dead|42|1|4|0
electrical short|15|29|8|3
fuel blockage|0|2|25|1
engine problem|1|0|8|12

```

Figure 15-2: a confusion matrix, for an expert system diagnosing situations in which a car won't start

Given the rule $a :- b, c$.

Assume a, b, c, d, e, f , and g are predicate expressions.

Assume b, c, d , and e are composed of one or more predicate expressions "and"ed and "or"d together.

```

Nature of/Predicate expressions/Suggested new rule
the new case/that are true
/in the new case (those
/not mentioned are false)

```

```

example|a and b and c|no change needed
example|a and b|a :- b.
example|a and b and c and d|no change needed
example|a and b and d|a :- b, e.
||where e is implied by both c and d
||(if no such e, try a:- b.)
example|f and b and c|g :- b, c.
||where g is implied by both a and f
||(if no such g, create a new rule f :- b, c.)
near miss|a and b and c|remove rule entirely
near miss|a and b|no change needed
near miss|a and b and c and d|a :- b, c, not(d).
near miss|a and b and d|no change needed
near miss|f and b and c|no change needed

```

Figure 15-3: suggested rule modifications given a new case

```

A/B/A AND B/A OR B/NOT A/
true|true|true|true|false
true|false|false|true|false
false|true|false|true|true
false|false|false|false|true

```

Figure A-1: summary of AND, OR, and NOT

Task: add 1000 to each salary in a set of 732
employee salaries, starting with employee #1

Task: add 100 to the Task: add 1000 to each salary in a set of 731
salary of employee #1 employee salaries, starting with employee #2

Task: add 1000 to the Task: add 1000 to each salary in a set of 730
salary of employee #2 employee salaries, starting with employee #3

Task: add 1000 to each salary in a set of 732
employee salaries, starting with employee #1

Task: add 1000 to each salary in Task: add 1000 to each salary in
a set of 366 employee salaries, a set of 366 employee salaries,
starting with employee #1 starting with employee #367

Task: add 1000 Task: add 1000 Task: add 1000 Task: add 1000
to each salary to each salary to each salary to each salary
in a set of 183
starting with starting with starting with starting with
employee #1 employee #184 employee #367 employee #550

Figure B-1: two different ways of doing an example recursion

Task: take the sum of 732 employee salaries, starting with employee #1

(the right box below calculates the result)

Task: take the sum of 731 employee salaries, starting with employee #2 to the salary of employee #1

(the right box below calculates the result)

Task: take the sum of 730 employee salaries, starting with employee #3 to the salary of employee #2

Figure B-2: an example of recursion of a function

array: list: stack using an array (insert and delete only at right end): queue using an array (insert at left end, delete at right end): tree: .sp 8

lattice: .sp 8

graph: .sp 8

Figure C-1: pictures of examples of some famous data structures

Built-in	Arguments	Can it succeed on	Description
<i>predicate//backtracking?</i>			
consult	a filename	no	loads a file
listing	a predicate name	no	prints database items with that predicate name
asserta	a fact	no	adds a fact to the front of the database
assertz	a fact	no	adds a fact to the rear of the database
retract	a fact	no	removes a fact
abolish	a predicate name	no	removes all facts with and its number the same predicate name of arguments
>	[infix] two numbers	no	greater-than
<	[infix] two numbers	no	less-than
=	[infix] two numbers	yes	equals
is	[infix] variable and	yes	arithmetic assignment a numeric expression
number	a variable	no	succeeds if argument is bound to a number
write	a variable or	no	prints on the terminal character string
read	a variable or	no	reads a word typed character string by the user
nl	none	no	issues carriage return to the terminal
not	a predicate	no	succeeds if querying its expression argument fails
fail	none	no	always fails
!	none	no	[see text]
var	a variable	no	succeeds if argument is not bound to anything
call	a predicate	yes	queries the predicate expression
clause	a rule left side	yes	binds text of a rule and rule right side to any variables
=..	[infix] a query	yes	converts a query to a and a list list or vice versa

Figure D-1: summary of built-in Prolog predicates assumed in this book

Standard Prolog/Micro-Prolog
feature/feature

```
consult(<filename>)|LOAD <filename>
listing(<predicatename>)|LIST <predicatename>
```

```

asserta(F) | (ADDCL F 0)
assertz(F) | (ADDCL F <number-of-F-facts-and-rules>)
retract(F) | (DELCL F)
abolish(F,2) | KILL F
?- |&.
:- | [see discussion]
, | [see discussion]
p(X); q(X) | (OR ((p X)) ((q X)))
not(p(X)) | (NOT p X)
X<Y | (LESS X Y)
X>Y | (LESS Y X)
X=3 | (EQ X 3)
X is 3 | (EQ X 3)
X is Y+Z | (SUM Y Z X)
X is Y-Z | (SUM X Z Y)
X is Y*Z | (TIMES Y Z X)
X is Y/Z | (TIMES X Z Y)
[a,b,X] | (a b X)
[X|Y] | (X | Y)
'Character string' | (Character string)
write(X) | (P X)
write(X),nl | (PP X)
read(X) | (R X)
clause(p(X),Y) | ((p X) | Y)
P=..L | [unnecessary--queries are lists]
call(p(X)) | (? ((p X))
fail | (FAIL)
!|!
number(X) | (NUM X)
var(X) | (VAR X)
/* comment */ /* comment
setof(X,p(X),X2) | (ISALL X2 X (p X))

```

Figure E-1: approximate equivalents between the Prolog assumed in this book and Micro-Prolog

```

          a
    b      c      d
e      f      g      h

```

Figure G-1: answer to problem 9-11.(a)

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